

## AT25DN011

### 1-Mbit, 2.3 V Minimum SPI Serial Flash Memory with Dual-Read Support

#### Features

- Single 2.3 V - 3.6 V Supply
- Serial Peripheral Interface (SPI) Compatible
  - Supports SPI Modes 0 and 3
  - Supports Dual Output Read
- 104 MHz Maximum Operating Frequency
  - Clock-to-Output ( $t_v$ ) of 6 ns
- Flexible, Optimized Erase Architecture for Code + Data Storage Applications
  - Small (256-Byte) Page Erase
  - Uniform 4-kByte Block Erase
  - Uniform 32-kByte Block Erase
  - Full Chip Erase
- Hardware Controlled Locking of Protected Sectors via  $\overline{WP}$  Pin
- 128-byte, One-Time Programmable (OTP) Security Register
- Flexible Programming
  - Byte/Page Program (1 to 256 Bytes)
- Fast Program and Erase Times
  - 1.25 ms Typical Page Program (256 Bytes) Time
  - 35 ms Typical 4-kByte Block Erase Time
  - 250 ms Typical 32-kByte Block Erase Time
- Automatic Checking and Reporting of Erase/Program Failures
- Software Controlled Reset
- JEDEC Standard Manufacturer and Device ID Read Methodology
- Low Power Dissipation
  - 350 nA Ultra-Deep Power-Down current (Typical)
  - 7.5  $\mu$ A Deep Power-Down Current (Typical)
  - 25  $\mu$ A Standby current (Typical)
  - 6 mA Active Read Current (Typical)
- Endurance: 100,000 Program/Erase Cycles
- Data Retention: 20 Years
- Complies with Full Industrial Temperature Range
- Industry Standard Green (Pb/Halide-free/RoHS Compliant) Package Options
  - 8-lead SOIC (150-mil)
  - 8-pad Ultra Thin DFN (2 x 3 x 0.6 mm)
  - 8-lead TSSOP Package

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## 1. Description

The AT25DN011 is a serial interface Flash memory device for use in a wide variety of high-volume consumer-based applications in which program code is shadowed from Flash memory into embedded or external RAM for execution. The flexible erase architecture of the AT25DN011, with its page erase granularity, is also ideal for data storage, eliminating the need for additional data storage devices.

The erase block sizes of the AT25DN011 have been optimized to meet the needs of today's code and data storage applications. By optimizing the size of the erase blocks, the memory space can be used much more efficiently. Because certain code modules and data storage segments must reside by themselves in their own erase regions, the wasted and unused memory space that occurs with large sectored and large block erase Flash memory devices can be greatly reduced. This increased memory space efficiency allows additional code routines and data storage segments to be added while still maintaining the same overall device density.

The device also contains a specialized One-Time Programmable (OTP) Security Register that can be used for unique device serialization, system-level Electronic Serial Number (ESN) storage, locked key storage, etc.

Designed for use in many different systems, the AT25DN011 supports read, program, and erase operations with a wide supply voltage range of 2.3 V to 3.6 V. No separate voltage is required for programming and erasing.

## 2. Pin Descriptions and Pinouts

Table 1. Pin Descriptions

Symbol	Name and Function	Asserted State	Type
$\overline{\text{CS}}$	<p><b>CHIP SELECT:</b> Asserting the <math>\overline{\text{CS}}</math> pin selects the device. When the <math>\overline{\text{CS}}</math> pin is deasserted, the device is deselected and normally placed in standby mode (not Deep Power-Down mode), and the SO pin is in a high-impedance state. When the device is deselected, data are not accepted on the SI pin.</p> <p>A high-to-low transition on the <math>\overline{\text{CS}}</math> pin is required to start an operation, and a low-to-high transition is required to end an operation. When ending an internally self-timed operation such as a program or erase cycle, the device does not enter the standby mode until the completion of the operation.</p>	Low	Input
SCK	<p><b>SERIAL CLOCK:</b> This pin is used to provide a clock to the device and is used to control the flow of data to and from the device. Command, address, and input data present on the SI pin is always latched in on the rising edge of SCK, while output data on the SO pin is always clocked out on the falling edge of SCK.</p>	-	Input
SI (I/O <sub>0</sub> )	<p><b>SERIAL INPUT:</b> The SI pin is used to shift data into the device. The SI pin is used for all data input including command and address sequences. Data on the SI pin is always latched in on the rising edge of SCK.</p> <p>With the Dual-Output Read commands, the SI Pin becomes an output pin (I/O<sub>0</sub>) in conjunction with other pins to allow two bits of data on (I/O<sub>1-0</sub>) to be clocked out on every falling edge of SCK. To maintain consistency with the SPI nomenclature, the SI (I/O<sub>0</sub>) pin is referenced as the SI pin unless specifically addressing the Dual-I/O modes, in which case it is referenced as I/O<sub>0</sub>. Data present on the SI pin are ignored whenever the device is deselected (<math>\overline{\text{CS}}</math> is deasserted).</p>	-	Input/Output
SO (I/O <sub>1</sub> )	<p><b>SERIAL OUTPUT:</b> The SO pin is used to shift data out from the device. Data on the SO pin is always clocked out on the falling edge of SCK.</p> <p>With the Dual-Output Read commands, the SO Pin remains an output pin (I/O<sub>1</sub>) in conjunction with other pins to allow two bits of data on (I/O<sub>1-0</sub>) to be clocked out on every falling edge of SCK. To maintain consistency with the SPI nomenclature, the SO (I/O<sub>1</sub>) pin is referenced as the SO pin unless specifically addressing the Dual-I/O modes, in which case it is referenced as I/O<sub>1</sub>. The SO pin is in a high-impedance state whenever the device is deselected (<math>\overline{\text{CS}}</math> is deasserted).</p>	-	Input/Output
$\overline{\text{WP}}$	<p><b>WRITE PROTECT:</b> The <math>\overline{\text{WP}}</math> pin controls the hardware locking feature of the device. See <a href="#">Section 9, Protection Commands and Features</a>, for more details on protection features and the <math>\overline{\text{WP}}</math> pin. The <math>\overline{\text{WP}}</math> pin is internally pulled-high and can be left floating if hardware controlled protection is not used. However, it is recommended that the <math>\overline{\text{WP}}</math> pin also be externally connected to V<sub>CC</sub> whenever possible.</p>	Low	Input
HOLD	<p><b>HOLD:</b> The HOLD pin is used to temporarily pause serial communication without deselecting or resetting the device. While the <math>\overline{\text{HOLD}}</math> pin is asserted, transitions on the SCK pin and data on the SI pin are ignored, and the SO pin is in a high-impedance state.</p> <p>The <math>\overline{\text{CS}}</math> pin must be asserted, and the SCK pin must be in the low state in order for a Hold condition to start. A Hold condition pauses serial communication only and does not have an effect on internally self-timed operations such as a program or erase cycle. See <a href="#">Section 12.7</a> for additional details on the Hold operation.</p> <p>The <math>\overline{\text{HOLD}}</math> pin is internally pulled-high and can be left floating if the Hold function is not used. However, it is recommended that the HOLD pin also be externally connected to V<sub>CC</sub> whenever possible.</p>	Low	Input
V <sub>CC</sub>	<p><b>DEVICE POWER SUPPLY:</b> The V<sub>CC</sub> pin is used to supply the source voltage to the device. Operations at invalid V<sub>CC</sub> voltages can produce spurious results; do not attempt this.</p>	-	Power
GND	<p><b>GROUND:</b> The ground reference for the power supply. Connect GND to the system ground.</p>	-	Power

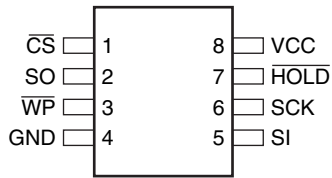


Figure 1. 8-SOIC Top View

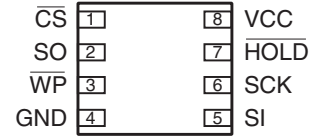


Figure 3. 8-UDFN (Top View)

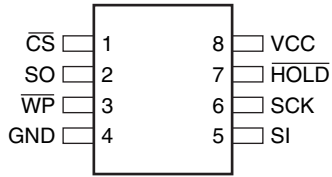


Figure 2. 8-TSSOP Top View



### 3. Block Diagram

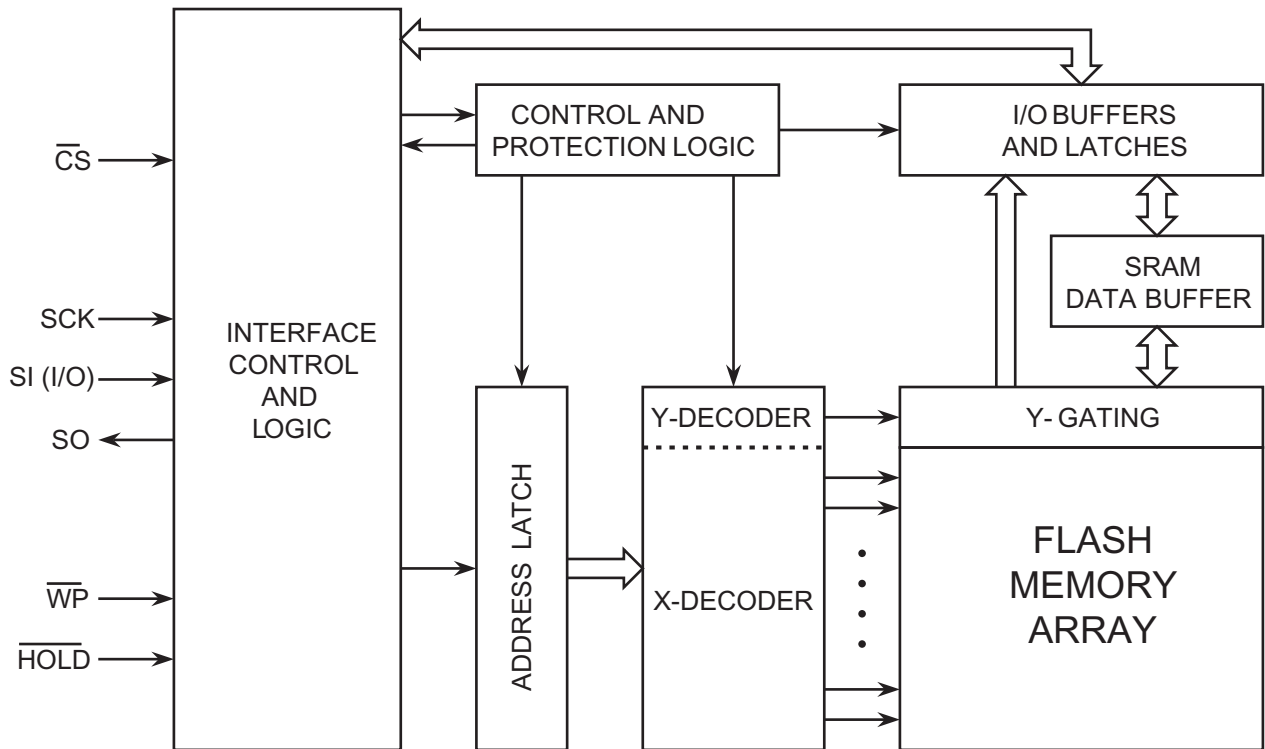


Figure 4. AT25DN011 Block Diagram

## 4. Memory Array

To provide the greatest flexibility, the memory array of the AT25DN011 can be erased in three levels of granularity, including a full-chip erase. The size of the erase blocks is optimized for both code and data storage applications, allowing both code and data segments to reside in their own erase regions. Figure 5 shows each erase level.

Block Erase Detail			Page Program Detail	
32KB Block Erase (52h Command)	4KB Block Erase (20h Command)	Block Address Range	1-256 Byte Page Program (02h Command)	Page Address Range
32KB	4KB	01FFFFh – 01F000h	256 Bytes	01FFFFh – 01FF00h
	4KB	01EFFFh – 01E000h	256 Bytes	01FEFFh – 01FE00h
	4KB	01DFFFh – 01D000h	256 Bytes	01DFFFh – 01FD00h
	4KB	01CFFFh – 01C000h	256 Bytes	01CFFFh – 01FC00h
	4KB	01BFFFh – 01B000h	256 Bytes	01BFFFh – 01FB00h
	4KB	01AFFFh – 01A000h	256 Bytes	01FAFFh – 01FA00h
	4KB	019FFFh – 019000h	256 Bytes	01F9FFh – 01F900h
	4KB	018FFFh – 018000h	256 Bytes	01F8FFh – 01F800h
32KB	4KB	017FFFh – 017000h	256 Bytes	01F7FFh – 01F700h
	4KB	016FFFh – 016000h	256 Bytes	01F6FFh – 01F600h
	4KB	015FFFh – 015000h	256 Bytes	01F5FFh – 01F500h
	4KB	014FFFh – 014000h	256 Bytes	01F4FFh – 01F400h
	4KB	013FFFh – 013000h	256 Bytes	01F3FFh – 01F300h
	4KB	012FFFh – 012000h	256 Bytes	01F2FFh – 01F200h
	4KB	011FFFh – 011000h	256 Bytes	01F1FFh – 01F100h
	4KB	010FFFh – 010000h	⋮	
32KB	4KB	00FFFFh – 00F000h	256 Bytes	000EFFh – 000E00h
	4KB	00EFFFh – 00E000h	256 Bytes	000DFFh – 000D00h
	4KB	00DFFFh – 00D000h	256 Bytes	000CFFh – 000C00h
	4KB	00CFFFh – 00C000h	256 Bytes	000BFFh – 000B00h
	4KB	00AFFFh – 00A000h	256 Bytes	000AFFh – 000A00h
	4KB	009FFFh – 009000h	256 Bytes	0009FFh – 000900h
	4KB	008FFFh – 008000h	256 Bytes	0008FFh – 000800h
	4KB	007FFFh – 007000h	256 Bytes	0007FFh – 000700h
32KB	4KB	006FFFh – 006000h	256 Bytes	0006FFh – 000600h
	4KB	005FFFh – 005000h	256 Bytes	0005FFh – 000500h
	4KB	004FFFh – 004000h	256 Bytes	0004FFh – 000400h
	4KB	003FFFh – 003000h	256 Bytes	0003FFh – 000300h
	4KB	002FFFh – 002000h	256 Bytes	0002FFh – 000200h
	4KB	001FFFh – 001000h	256 Bytes	0001FFh – 000100h
	4KB	000FFFh – 000000h	256 Bytes	0000FFh – 000000h

Figure 5. Memory Architecture Diagram

## 5. Device Operation

The AT25DN011 is controlled by a set of commands that are sent from a host controller, commonly referred to as the SPI Master. The SPI Master communicates with the AT25DN011 through the SPI bus, which consists of four signal lines: Chip Select ( $\overline{CS}$ ), Serial Clock (SCK), Serial Input (SI), and Serial Output (SO).

The SPI protocol defines four modes of operation (0, 1, 2, or 3), with each mode differing in respect to the SCK polarity and phase, as well as how the polarity and phase control the flow of data on the SPI bus. The AT25DN011 supports the two most common modes, SPI Modes 0 and 3. The only difference between SPI Modes 0 and 3 is the polarity of the SCK signal when in the inactive state (when the SPI Master is in standby mode and not transferring any data). With SPI Modes 0 and 3, data is latched in on the rising edge of SCK and output on the falling edge of SCK.

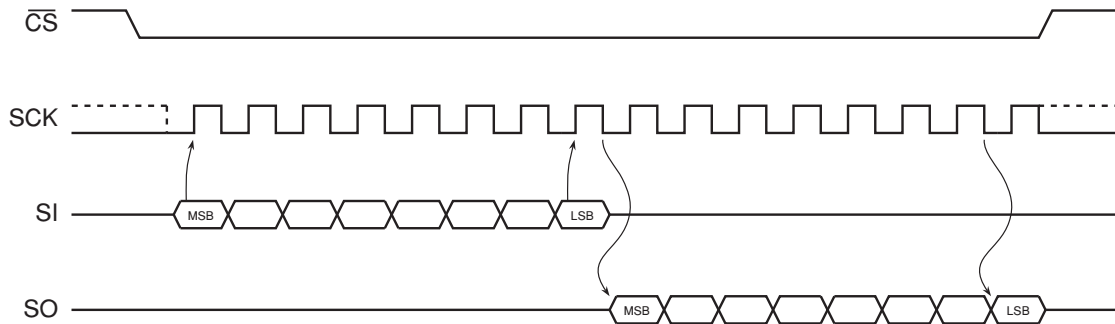


Figure 6. SPI Mode 0 and 3

### 5.1 Dual Output Read

The AT25DN011 features a Dual-Output Read mode that allow two bits of data to be clocked out of the device every clock cycle to improve throughput. To do this, both the SI and SO pins are used as outputs for the transfer of data bytes. With the Dual-Output Read Array command, the SI pin becomes an output along with the SO pin.

## 6. Commands and Addressing

A valid command or operation must be started by first asserting the  $\overline{\text{CS}}$  pin. After the  $\overline{\text{CS}}$  pin has been asserted, the host controller must then clock out a valid 8-bit opcode on the SPI bus. Following the opcode, command-dependent information, such as address and data bytes, is then clocked out by the host controller. All opcode, address, and data bytes are transferred with the most-significant bit (MSB) first. An operation is ended by deasserting the  $\overline{\text{CS}}$  pin.

Opcodes not supported by the AT25DN011 are ignored by the device, and no operation is started. The device continues to ignore any data presented on the SI pin until the start of the next operation ( $\overline{\text{CS}}$  pin being deasserted and then reasserted). Also, if the  $\overline{\text{CS}}$  pin is deasserted before complete opcode and address information is sent to the device, then no operation is performed, and the device returns to the idle state and waits for the next operation.

Addressing of the device requires a total of three bytes of information to be sent; these represent address bits A23-A0. Since the upper address limit of the AT25DN011 memory array is 01FFFFh, address bits A23-A17 are always ignored by the device.

**Table 2. Command Listing**

Command	Opcode	Clock Frequency	Address Bytes	Dummy Bytes	Data Bytes	
<b>Read Commands</b>						
Read Array	0Bh	0000 1011	Up to 104 MHz	3	1	1+
	03h	0000 0011	Up to 33 MHz <sup>1</sup>	3	0	1+
Dual Output Read	3Bh	0011 1011	Up to 50 MHz	3	1	1+
<b>Program and Erase Commands</b>						
Page Erase	81h	1000 0001	Up to 104 MHz	3	0	0
Block Erase (4 kBytes)	20h	0010 0000	Up to 104 MHz	3	0	0
Block Erase (32 kBytes)	52h	0101 0010	Up to 104 MHz	3	0	0
	D8h	1101 1000	Up to 104 MHz	3	0	0
Chip Erase	60h	0110 0000	Up to 104 MHz	0	0	0
	C7h	1100 0111	Up to 104 MHz	0	0	0
Chip Erase (Legacy Command)	62h	0110 0010	Up to 104 MHz	0	0	0
Byte/Page Program (1 to 256 Bytes)	02h	0000 0010	Up to 104 MHz	3	0	1+
<b>Protection Commands</b>						
Write Enable	06h	0000 0110	Up to 104 MHz	0	0	0
Write Disable	04h	0000 0100	Up to 104 MHz	0	0	0
<b>Security Commands</b>						
Program OTP Security Register	9Bh	1001 1011	Up to 104 MHz	3	0	1+
Read OTP Security Register	77h	0111 0111	Up to 104 MHz	3	2	1+
<b>Status Register Commands</b>						
Read Status Register	05h	0000 0101	Up to 104 MHz	0	0	1+
Write Status Register Byte 1	01h	0000 0001	Up to 104 MHz	0	0	1
Write Status Register Byte 2	31h	0011 0001	Up to 104 MHz	0	0	1
<b>Miscellaneous Commands</b>						
Reset	F0h	1111 0000	Up to 104 MHz	0	0	1 (D0h)
Read Manufacturer and Device ID	9Fh	1001 1111	Up to 104 MHz	0	0	1 to 4
Read ID (Legacy Command)	15h	0001 0101	Up to 104 MHz	0	0	2
Deep Power-Down	B9h	1011 1001	Up to 104 MHz	0	0	0
Resume from Deep Power-Down	ABh	1010 1011	Up to 104 MHz	0	0	0
Ultra-Deep Power-Down	79h	0111 1001	Up to 104 MHz	0	0	0

## 7. Read Commands

### 7.1 Read Array

The Read Array command sequentially reads a continuous stream of data from the device by providing the clock signal once the initial starting address is specified. The device incorporates an internal address counter that automatically increments every clock cycle.

Two opcodes (0Bh and 03h) can be used for the Read Array command. The use of each opcode depends on the maximum clock frequency used to read data from the device. The 0Bh opcode can be used at any clock frequency up to the maximum specified by  $f_{CLK}$ , and the 03h opcode can be used for lower frequency read operations up to the maximum specified by  $f_{RDLF}$ .

To perform the Read Array operation, the  $\overline{CS}$  pin must first be asserted and the appropriate opcode (0Bh or 03h) must be clocked into the device. After the opcode has been clocked in, the three address bytes must be clocked in to specify the starting address location of the first byte to read within the memory array. Following the three address bytes, an additional dummy byte must be clocked into the device if the 0Bh opcode is used for the Read Array operation.

After the three address bytes (and the dummy byte if using opcode 0Bh) have been clocked in, additional clock cycles result in data being output on the SO pin. The data is always output with the MSB of a byte first. When the last byte (01FFFFh) of the memory array has been read, the device continues reading at the beginning of the array (000000h). No delays are incurred when wrapping around from the end of the data array to the beginning of the array.

Deasserting the  $\overline{CS}$  pin stops the read operation and puts the SO pin into high-impedance state. The  $\overline{CS}$  pin can be deasserted at any time and does not require a full byte of data be read.

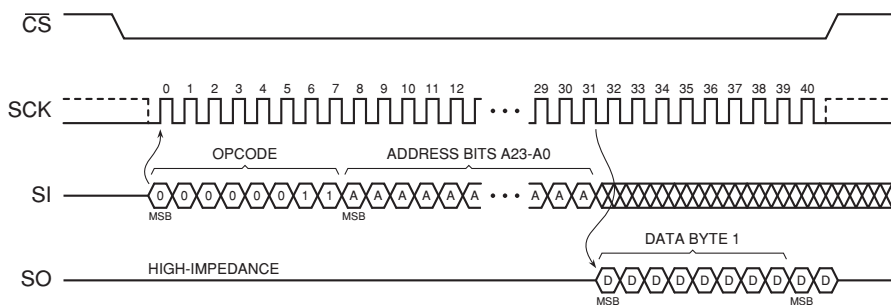


Figure 7. Read Array - 03h Opcode

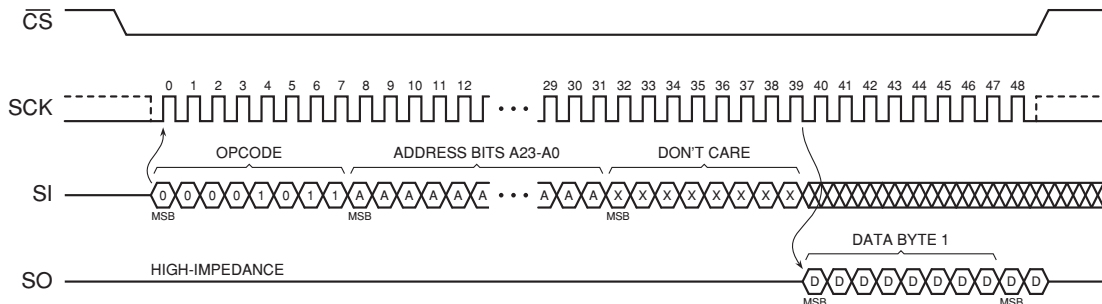


Figure 8. Read Array - 0Bh Opcode



## 8. Program and Erase Commands

### 8.1 Byte/Page Program

The Byte/Page Program command allows anywhere from a single byte of data to 256 bytes of data to be programmed into previously erased memory locations. An erased memory location is one that has all eight bits set to the logical 1 state (a byte value of FFh). Before a Byte/Page Program command can be started, the Write Enable command must have been previously issued to the device (see [Section 9.1, Write Enable](#)) to set the Write Enable Latch (WEL) bit of the Status Register to a logical 1 state.

To perform a Byte/Page Program command, an opcode of 02h must be clocked into the device, followed by the three address bytes denoting the first byte location of the memory array to begin programming at. After the address bytes have been clocked in, data can then be clocked into the device and stored in an internal buffer.

If the starting memory address denoted by A23-A0 does not fall on an even 256-byte page boundary (A7-A0 are not all 0), then special circumstances regarding which memory locations to be programmed apply. In this situation, any data that are sent to the device that goes beyond the end of the page wrap around to the beginning of the same page. For example, if the starting address denoted by A23-A0 is 0000FEh, and three bytes of data are sent to the device, then the first two bytes of data are programmed at addresses 0000FEh and 0000FFh, while the last byte of data is programmed at address 000000h. The remaining bytes in the page (addresses 000001h through 0000FDh) are not programmed and remain in the erased state (FFh). Also, if more than 256 bytes of data are sent to the device, then only the last 256 bytes sent are latched into the internal buffer.

When the  $\overline{CS}$  pin is deasserted, the device takes the data stored in the internal buffer and programs it into the appropriate memory array locations based on the starting address specified by A23-A0 and the number of data bytes sent to the device. If less than 256 bytes of data were sent to the device, then the remaining bytes within the page are not programmed and remain in the erased state (FFh). The programming of the data bytes is internally self-timed and takes place in a time of  $t_{PP}$  or  $t_{BP}$  if only programming a single byte.

The three address bytes and at least one complete byte of data must be clocked into the device before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on even byte boundaries (multiples of eight bits); otherwise, the device stops the operation, and no data are programmed into the memory array. Also, if the memory is in the protected state (see [Section 9.3, Block Protection](#)), then the Byte/Page Program command is not executed, and the device returns to the idle state once the  $\overline{CS}$  pin has been deasserted. The WEL bit in the Status Register is reset to the logical 0 state if the program cycle stops due to an incomplete address being sent, an incomplete byte of data being sent, the  $\overline{CS}$  pin being deasserted on uneven byte boundaries, or because the memory location to be programmed is protected.

While the device is programming, the Status Register can be read and indicates that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the  $t_{BP}$  or  $t_{PP}$  time to determine if the data bytes have finished programming. At some point before the program cycle completes, the WEL bit in the Status Register is reset to the logical 0 state.

The device also incorporates an intelligent programming algorithm that can detect when a byte location fails to program properly. If a programming error arises, it is indicated by the EPE bit in the Status Register.

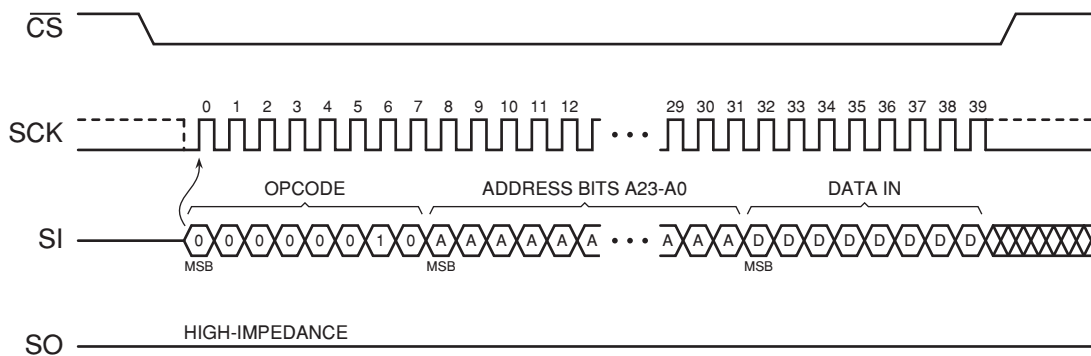


Figure 10. Byte Program

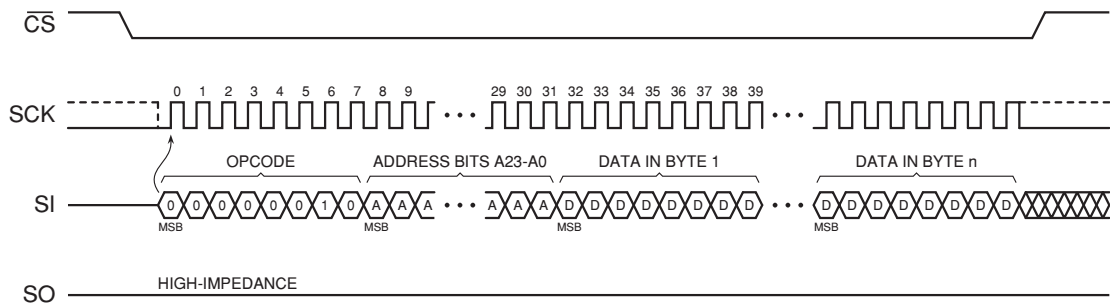


Figure 11. Page Program

## 8.2 Page Erase

Page Erase for 1 Mbit, 512 Pages [nine (9) page address bits, PA<8:0>] of 256 bytes each.

The Page Erase command can be used to individually erase any page in the main memory array. The Main Memory Byte/Page Program command can be used at a later time.

To perform a Page Erase with the standard page size (256 bytes), an opcode of 81h must be clocked into the device followed by three address bytes comprised of:

Byte 0: 81h the page erase command code

Byte 1: XXXX XXX, PA8; which is seven (7) dummy bits, and one (1) page address bit PA8

Byte 2: PA<7:0>; which is eight (8) page address bits

Byte 3: XXXX XXXX; which is eight (8) dummy bits

When a low-to-high transition occurs on the  $\overline{CS}$  pin, the device erases the selected page (the erased state is a Logic 1). The erase operation is internally self-timed and takes place in a maximum time of  $t_{PE}$ . During this time, the  $\overline{RDY}/BUSY$  bit in the Status Register indicates that the device is busy.

The device also incorporates an intelligent erase algorithm that can detect when a byte location fails to erase properly. If an erase error arises, it is indicated by the EPE bit in the Status Register.



### 8.3 Block Erase

A block of 4 or 32 kBytes can be erased (all bits set to the logical 1 state) in a single operation by using one of three different opcodes for the Block Erase command. An opcode of 20h is used for a 4-kByte erase, and an opcode of 52h or D8h is used for a 32-kByte erase. Before a Block Erase command can be started, the Write Enable command must have been previously issued to the device to set the WEL bit of the Status Register to a logical 1 state.

To perform a Block Erase, the  $\overline{CS}$  pin must first be asserted and the appropriate opcode (20h, 52h, or D8h) must be clocked into the device. After the opcode has been clocked in, the three address bytes specifying an address within the 4- or 32-kByte block to be erased must be clocked in. Any additional data clocked into the device are ignored. When the  $\overline{CS}$  pin is deasserted, the device erases the appropriate block. The erasing of the block is internally self-timed and takes place in a time of  $t_{BLKE}$ .

Since the Block Erase command erases a region of bytes, the lower order address bits do not need to be decoded by the device. Thus, for a 4-kByte erase, address bits A11-A0 are ignored by the device, and their values can be either a logical 1 or 0. For a 32-kByte erase, address bits A14-A0 are ignored by the device. Despite the lower order address bits not being decoded by the device, the complete three address bytes must still be clocked into the device before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, the device stops the operation, and no erase operation is performed.

If the memory is in the protected state, then the Block Erase command is not executed, and the device returns to the idle state once the  $\overline{CS}$  pin has been deasserted.

The WEL bit in the Status Register is reset to the logical 0 state if the erase cycle stops due to an incomplete address being sent, the  $\overline{CS}$  pin being deasserted on uneven byte boundaries, or because a memory location within the region to be erased is protected.

While the device is executing a successful erase cycle, the Status Register can be read and indicates that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the  $t_{BLKE}$  time to determine if the device has finished erasing. At some point before the erase cycle completes, the WEL bit in the Status Register is reset to the logical 0 state.

The device also incorporates an intelligent erase algorithm that can detect when a byte location fails to erase properly. If an erase error occurs, it is indicated by the EPE bit in the Status Register.

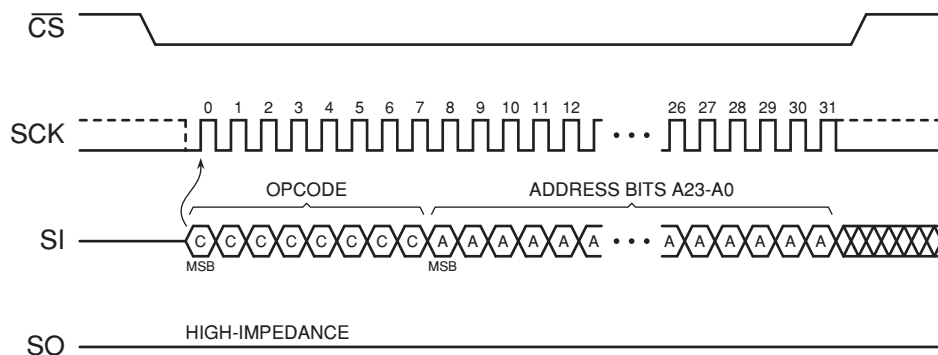


Figure 12. Block Erase

## 8.4 Chip Erase

The entire memory array can be erased in a single operation by using the Chip Erase command. Before a Chip Erase command can be started, the Write Enable command must have been previously issued to the device to set the WEL bit of the Status Register to a logical 1 state.

Three opcodes (60h, 62h, and C7h) can be used for the Chip Erase command. There is no difference in device functionality when using the three opcodes, so they can be used interchangeably. To perform a Chip Erase, one of the three opcodes must be clocked into the device. Since the entire memory array is to be erased, no address bytes need to be clocked into the device, and any data clocked in after the opcode are ignored. When the  $\overline{CS}$  pin is deasserted, the device erases the entire memory array. The erasing of the device is internally self-timed and takes place in a time of  $t_{CHPE}$ .

The complete opcode must be clocked into the device before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, no erase is performed. Also, if the memory array is in the protected state, then the Chip Erase command is not executed, and the device returns to the idle state once the  $\overline{CS}$  pin has been deasserted. The WEL bit in the Status Register is reset to the logical 0 state if the  $\overline{CS}$  pin is deasserted on uneven byte boundaries or if the memory is in the protected state.

While the device is executing a successful erase cycle, the Status Register can be read and indicates that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the  $t_{CHPE}$  time to determine if the device has finished erasing. At some point before the erase cycle completes, the WEL bit in the Status Register is reset to the logical 0 state.

The device also incorporates an intelligent erase algorithm that can detect when a byte location fails to erase properly. If an erase error occurs, it is indicated by the EPE bit in the Status Register.

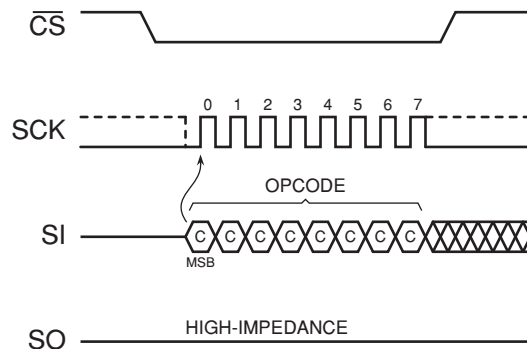


Figure 13. Chip Erase

## 9. Protection Commands and Features

### 9.1 Write Enable

The Write Enable command is used to set the Write Enable Latch (WEL) bit in the Status Register to a logical 1 state. The WEL bit must be set before a Byte/Page Program, erase, Program OTP Security Register, or Write Status Register command can be executed. This makes the issuance of these commands a two step process, thereby reducing the chances of a command being accidentally or erroneously executed. If the WEL bit in the Status Register is not set prior to the issuance of one of these commands, then the command is not executed.

To issue the Write Enable command, the  $\overline{CS}$  pin must first be asserted and the opcode of 06h must be clocked into the device. No address bytes need to be clocked into the device, and any data clocked in after the opcode are ignored. When the  $\overline{CS}$  pin is deasserted, the WEL bit in the Status Register is set to a logical 1. The complete opcode must be clocked into the device before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, the device stops the operation, and the state of the WEL bit does not change.

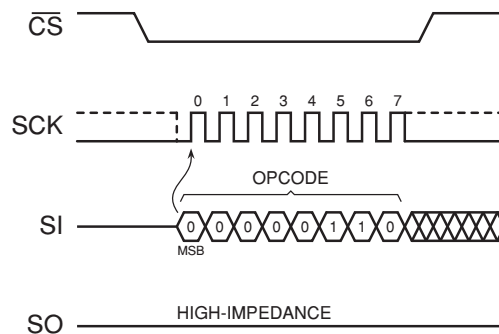


Figure 14. Write Enable

### 9.2 Write Disable

The Write Disable command is used to reset the Write Enable Latch (WEL) bit in the Status Register to the logical 0 state. With the WEL bit reset, all Byte/Page Program, erase, Program OTP Security Register, and Write Status Register commands are not executed. Other conditions can also cause the WEL bit to be reset; for more details, refer to the WEL bit section of the Status Register description.

To issue the Write Disable command, the  $\overline{CS}$  pin must first be asserted and the opcode of 04h must be clocked into the device. No address bytes need to be clocked into the device, and any data clocked in after the opcode are ignored. When the  $\overline{CS}$  pin is deasserted, the WEL bit in the Status Register is reset to a logical 0. The complete opcode must be clocked into the device before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, the device stops the operation, and the state of the WEL bit does not change.

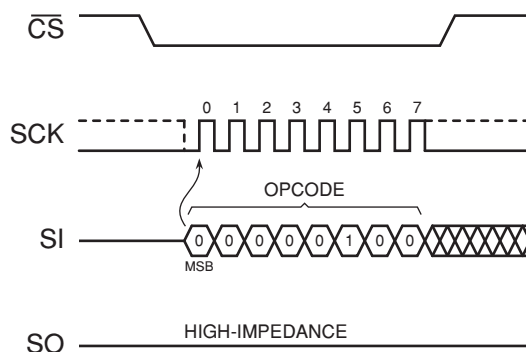


Figure 15. Write Disable

## 9.3 Block Protection

The device can be software protected against erroneous or malicious program or erase operations by using the Block Protection feature. Using the Write Status Register command to change the value of the Block Protection (BP0) bit in the Status Register. [Table 3](#) outlines the two states of the BP0 bit and the associated protection area.

**Table 3. Memory Array Protection**

Protection Level	BP0	Protected Memory Address
None	0	None
Full Memory	1	00000h - 01FFFFh

When the BP0 bit of the Status Register is in the logical 1 state, the entire memory array is protected against program or erase operations. Any attempts to send a Byte/Page Program command, a Block Erase command, or a Chip Erase command are ignored by the device.

As a safeguard against accidental or erroneous protecting or unprotecting of the memory array, the BP0 bit itself can be locked from updates by using the  $\overline{WP}$  pin and the BPL (Block Protection Locked) bit of the Status Register (see [Section 9.4, Protected States and the Write Protect Pin](#), for more details).

The BP0 bit of the Status Register is a nonvolatile bit; thus, the BP0 bit retains its state even after the device has been power-cycled. Ensure that BP0 is in the logical 1 state before powering down for those applications that want the memory array fully protected on power-up. The default state for BP0 when shipped is 0.

## 9.4 Protected States and the Write Protect Pin

The  $\overline{WP}$  pin is not linked to the memory array itself and has no direct effect on the protection status of the memory array. Instead, the  $\overline{WP}$  pin, with the BPL (Block Protection Locked) bit in the Status Register, is used to control the hardware-locking mechanism of the device. For hardware locking to be active, two conditions must be met: the  $\overline{WP}$  pin must be asserted, and the BPL bit must be in the logical 1 state.

When hardware locking is active, the Block Protection (BP0) bit and the BPL bit are locked. Thus, if the memory array is protected, it is locked in the protected state; if the memory array is unprotected, it is locked in the unprotected state. These states cannot be changed as long as hardware locking is active, so the Write Status Register command is ignored. To modify the protection status of the memory array, first deassert the  $\overline{WP}$  pin, then reset the BPL bit in the Status Register to the logical 0 state using the Write Status Register command.

If the  $\overline{WP}$  pin is permanently connected to GND, then once the BPL bit is set to a logical 1, the only way to reset the bit to the logical 0 state is to power-cycle the device. This allows a system to power-up with all sectors software protected but not hardware locked. Thus, sectors can be unprotected and protected as needed and then hardware locked at a later time by simply setting the BPL bit in the Status Register.

When the  $\overline{WP}$  pin is deasserted, or if the  $\overline{WP}$  pin is permanently connected to  $V_{CC}$ , the BPL bit in the Status Register can be set to a logical 1, but doing so does not lock the BP0 bit.

[Table 4](#) lists the various protection and locking states of the device.

**Table 4. Hardware and Software Locking**

$\overline{WP}$	BPL	Locking	BPL Change Allowed	BP0 and Protection Status
0	0		Can be modified from 0 to 1	BP0 bit unlocked and modifiable using the Write Status Register command. Memory array can be protected and unprotected freely.
0	1	Hardware Locked	Locked	BP0 bit locked in current state. The Write Status Register command has no affect. Memory array is locked in current protected or unprotected state.
1	0		Can be modified from 0 to 1	BP0 bit unlocked and modifiable using the Write Status Register command. Memory array can be protected and unprotected freely.
1	1		Can be modified from 1 to 0	BP0 bit unlocked and modifiable using the Write Status Register command. Memory array can be protected and unprotected freely.

## 10. Security Commands

### 10.1 Program OTP Security Register

The device contains a specialized OTP (One-Time Programmable) Security Register that can be used for purposes such as unique device serialization, system-level Electronic Serial Number (ESN) storage, locked key storage, etc. The OTP Security Register is independent of the main Flash memory array and is comprised of a total of 128 bytes of memory divided into two portions. The first 64 bytes (byte locations 0 through 63) of the OTP Security Register are allocated as a one-time user-programmable space. Once these 64 bytes have been programmed, they cannot be erased or reprogrammed. The remaining 64 bytes of the OTP Security Register (byte locations 64 through 127) are factory programmed by Renesas Electronics and contain a unique value for each device. The factory programmed data is fixed and cannot be changed.

**Table 5. OTP Security Register**

Security Register Byte Number									
0	1	...	62	63	64	65	...	126	127
One-Time User Programmable					Factory Programmed by Renesas Electronics				

The user-programmable portion of the OTP Security Register does not need to be erased before it is programmed. Also, the Program OTP Security Register command operates on the entire 64-byte user-programmable portion of the OTP Security Register at one time. Once the user-programmable space has been programmed with any number of bytes, the user-programmable space cannot be programmed again; thus, it is not possible to only program the first two bytes of the register and then program the remaining 62 bytes at a later time.

Before the Program OTP Security Register command can be issued, the Write Enable command must have been previously issued to set the WEL bit in the Status Register to a logical 1. To program the OTP Security Register, the  $\overline{CS}$  pin must first be asserted and an opcode of 9Bh must be clocked into the device followed by the three address bytes denoting the first byte location of the OTP Security Register to begin programming at. Since the size of the user-programmable portion of the OTP Security Register is 64 bytes, the upper order address bits do not need to be decoded by the device. Thus, address bits A23-A6 are ignored by the device, and their values can be either a logical 1 or 0. After the address bytes have been clocked in, data can then be clocked into the device and stored in the internal buffer.

If the starting memory address denoted by A23-A0 does not start at the beginning of the OTP Security Register memory space (A5-A0 are not all 0), then special circumstances regarding which OTP Security Register locations to be programmed apply. In this situation, any data sent to the device that goes beyond the end of the 64-byte user-programmable space wrap around to the beginning of the OTP Security Register. For example, if the starting address denoted by A23-A0 is 00003Eh, and three bytes of data are sent to the device, then the first two bytes of data are programmed at OTP Security Register addresses 00003Eh and 00003Fh, while the last byte of data is programmed at address 000000h. The remaining bytes in the OTP Security Register (addresses 000001h through 00003Dh) are not programmed and remain in the erased state (FFh). Also, if more than 64 bytes of data are sent to the device, then only the last 64 bytes sent are latched into the internal buffer.

When the  $\overline{CS}$  pin is deasserted, the device takes the data stored in the internal buffer and program it into the appropriate OTP Security Register locations based on the starting address specified by A23-A0 and the number of data bytes sent to the device. If less than 64 bytes of data were sent to the device, then the remaining bytes within the OTP Security Register are not programmed and remain in the erased state (FFh). The programming of the data bytes is internally self-timed and takes place in a time of  $t_{OTP}$ .

The three address bytes and at least one complete byte of data must be clocked into the device before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on even byte boundaries (multiples of eight bits); otherwise, the device stops the operation, and the user-programmable portion of the OTP Security Register is not programmed. The WEL bit in the Status Register is reset to the logical 0 state if the OTP Security Register program cycle stops due to an incomplete address being sent, an incomplete byte of data being sent, the  $\overline{CS}$  pin being deasserted on

uneven byte boundaries, or because the user-programmable portion of the OTP Security Register was previously programmed.

While the device is programming the OTP Security Register, the Status Register can be read and indicates that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the  $t_{OTPP}$  time to determine if the data bytes have finished programming. At some point before the OTP Security Register programming completes, the WEL bit in the Status Register is reset to the logical 0 state.

If the device is powered-down during the OTP Security Register program cycle, then the contents of the 64-byte user programmable portion of the OTP Security Register cannot be guaranteed and cannot be programmed again.

The Program OTP Security Register command uses the internal 256-buffer for processing. Thus, the contents of the buffer are altered from the previous state when this command is issued.

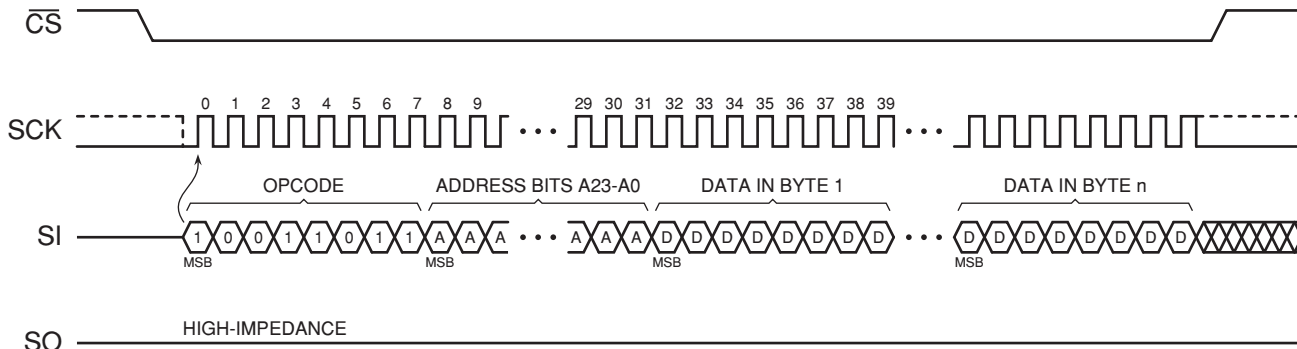


Figure 16. Program OTP Security Register

### 10.2 Read OTP Security Register

The OTP Security Register can be sequentially read in a similar fashion to the Read Array operation up to the maximum clock frequency specified by  $f_{CLK}$ . To read the OTP Security Register, the  $\overline{CS}$  pin must first be asserted and the opcode of 77h must be clocked into the device. After the opcode has been clocked in, the three address bytes must be clocked in to specify the starting address location of the first byte to read within the OTP Security Register. Following the three address bytes, two dummy bytes must be clocked into the device before data can be output.

After the three address bytes and the dummy bytes have been clocked in, additional clock cycles result in OTP Security Register data being output on the SO pin. When the last byte (00007Fh) of the OTP Security Register has been read, the device continues reading at the beginning of the register (000000h). No delays are incurred when wrapping around from the end of the register to the beginning of the register.

Deasserting the  $\overline{CS}$  pin stops the read operation and puts the SO pin into a high-impedance state. The  $\overline{CS}$  pin can be deasserted at any time and does not require that a full byte of data be read.

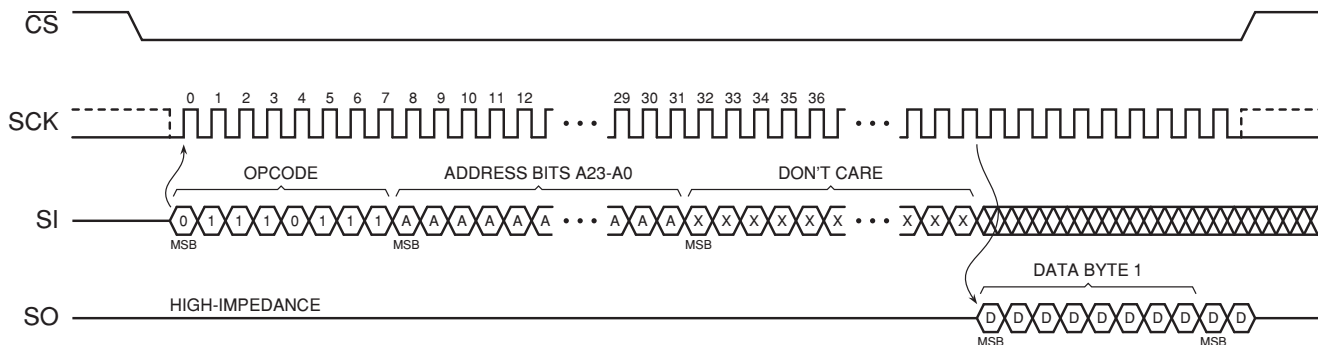


Figure 17. Read OTP Security Register

## 11. Status Register Commands

### 11.1 Read Status Register

The Status Register can be read to determine the device's ready/busy status, as well as the status of many other functions such as Hardware Locking and Block Protection. The Status Register can be read at any time, including during an internally self-timed program or erase operation. The Status Register consists of two bytes.

To read the Status Register, the  $\overline{CS}$  pin must first be asserted and the opcode of 05h must be clocked into the device. After the opcode has been clocked in, the device begins outputting Status Register data on the SO pin during every subsequent clock cycle. After the last bit (bit 0) of Status Register Byte 1 has been clocked out, the first bit (bit 7) of Status Register Byte 2 is clocked out. After the last bit (bit 0) of Status Register Byte 2 has been clocked out, the sequence repeats itself, starting again with bit 7 of Status Register Byte 1, as long as the  $\overline{CS}$  pin remains asserted and the clock pin is being pulsed. The data in the Status Register is constantly being updated, so each repeating sequence outputs new data.

Deasserting the  $\overline{CS}$  pin stops the Read Status Register operation and puts the SO pin into a high-impedance state. The  $\overline{CS}$  pin can be deasserted at any time and does not require that a full byte of data be read.

Table 6. Status Register Format

Bit <sup>1</sup>	Name		Type <sup>2</sup>	Description	
7	BPL	Block Protection Locked	R/W	0	BP0 bit unlocked (default).
				1	BP0 bit locked in current state when $\overline{WP}$ asserted.
6	RES	Reserved for future use	R	0	Reserved for future use.
5	EPE	Erase/Program Error	R	0	Erase or program operation was successful.
				1	Erase or program error detected.
4	WPP	Write Protect ( $\overline{WP}$ ) Pin Status	R	0	$\overline{WP}$ is asserted.
				1	$\overline{WP}$ is deasserted.
3	RES	Reserved for future use	R	0	Reserved for future use.
2	BP0	Block Protection	R/W	0	Entire memory array is unprotected.
				1	Entire memory array is protected.
1	WEL	Write Enable Latch Status	R	0	Device is not write enabled (default).
				1	Device is write enabled.
0	$\overline{RDY/BSY}$	Ready/Busy Status	R	0	Device is ready.
				1	Device is busy with an internal operation.

1. Only bits 7 and 2 of the Status Register can be modified when using the Write Status Register command.

2. R/W = Readable and writable; R = Readable only.

#### 11.1.1 BPL Bit

The BPL bit is used to control whether the Block Protection (BP0) bit can be modified or not. When the BPL bit is in the logical 1 state and the  $\overline{WP}$  pin is asserted, the BP0 bit is locked and cannot be modified. The memory array is locked in the current protected or unprotected state.

When the BPL bit is in the logical 0 state, the BP0 bit is unlocked and can be modified. The BPL bit defaults to the logical 0 state after device power-up.

The BPL bit can be modified freely whenever the  $\overline{WP}$  pin is deasserted. However, if the  $\overline{WP}$  pin is asserted, then the BPL bit can only be changed from a logical 0 (BP0 bit unlocked) to a logical 1 (BP0 bit locked). In order to reset the BPL bit to a logical 0 using the Write Status Register command, the  $\overline{WP}$  pin first must be deasserted.

The BPL and BP0 bits are the only bits of the Status Register that can be user modified via the Write Status Register command.

### 11.1.2 EPE Bit

The EPE bit indicates whether the last erase or program operation completed successfully or not. If at least one byte during the erase or program operation did not erase or program properly, then the EPE bit is set to the logical 1 state. The EPE bit is not set if an erase or program operation stops for any reason such as an attempt to erase or program the memory when it is protected or if the WEL bit is not set prior to an erase or program operation. The EPE bit is updated after every erase and program operation.

### 11.1.3 WPP Bit

The WPP bit can be read to determine if the  $\overline{WP}$  pin has been asserted or not.

### 11.1.4 BP0 Bit

The BP0 bits provides feedback on the software protection status for the device. Also, the BP0 bit can also be modified to change the state of the software protection to allow the entire memory array to be protected or unprotected. When the BP0 bit is in the logical 0 state, then the entire memory array is unprotected. When the BP0 bit is in the logical 1 state, then the entire memory array is protected against program and erase operations.

### 11.1.5 WEL Bit

The WEL bit indicates the current status of the internal Write Enable Latch. When the WEL bit is in the logical 0 state, the device does not accept any Byte/Page Program, erase, Program OTP Security Register, or Write Status Register commands. The WEL bit defaults to the logical 0 state after a device power-up or reset operation. Also, the WEL bit is reset to the logical 0 state automatically under the following conditions:

- Write Disable operation completes successfully
- Write Status Register operation completes successfully or stops
- Program OTP Security Register operation completes successfully or stops
- Byte/Page Program operation completes successfully or stops
- Block Erase operation completes successfully or stops
- Chip Erase operation completes successfully or stops
- Hold condition stops

If the WEL bit is in the logical 1 state, it is not reset to a logical 0 if an operation stops due to an incomplete or unrecognized opcode being clocked into the device before the  $\overline{CS}$  pin is deasserted. In order for the WEL bit to be reset when an operation stops prematurely, the entire opcode for a Byte/Page Program, erase, Program OTP Security Register, or Write Status Register command must have been clocked into the device.

### 11.1.6 $\overline{RDY/BSY}$ Bit

The  $\overline{RDY/BSY}$  bit is used to determine whether or not an internal operation, such as a program or erase, is in progress. To poll the  $\overline{RDY/BSY}$  bit to detect the completion of a program or erase cycle, new Status Register data must be continually clocked out of the device until the state of the  $\overline{RDY/BSY}$  bit changes from a logical 1 to a logical 0. Note that the  $\overline{RDY/BSY}$  bit can be read either from Status Register Byte 1 or from Status Register Byte 2.



### 11.1.7 RSTE Bit

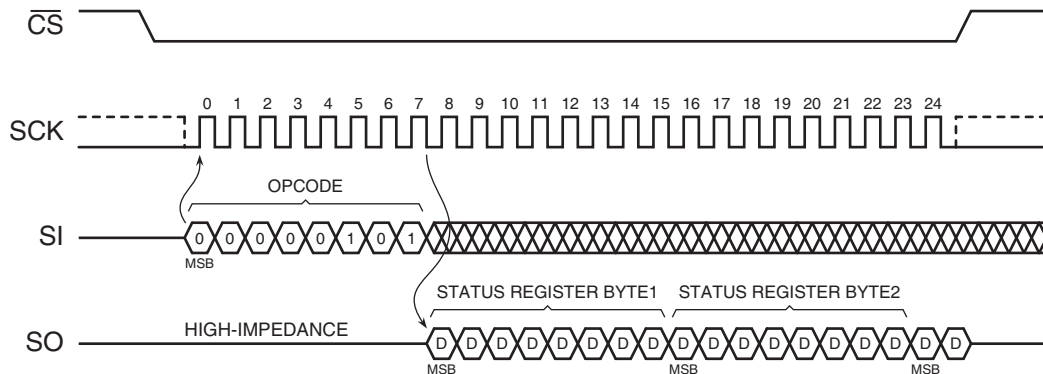
The RSTE bit is used to enable or disable the Reset command. When the RSTE bit is in the Logical 0 state (the default state after power-up), the Reset command is disabled and any attempts to reset the device using the Reset command are ignored. When the RSTE bit is in the Logical 1 state, the Reset command is enabled.

The RSTE bit retains its state as long as power is applied to the device. Once set to the Logical 1 state, the RSTE bit remains in that state until it is modified using the Write Status Register Byte 2 command or until the device has been power cycled. The Reset command itself does not change the state of the RSTE bit.

**Table 7. Status Register Format – Byte 2**

Bit <sup>(1)</sup>	Name		Type	Description	
7	RES	Reserved for future use	R	0	Reserved for future use
6	RES	Reserved for future use	R	0	Reserved for future use
5	RES	Reserved for future use	R	0	Reserved for future use
4	RSTE	Reset Enabled	R/W	0	Reset command is disabled (default)
				1	Reset command is enabled
3	RES	Reserved for future use	R	0	Reserved for future use
2	RES	Reserved for future use	R	0	Reserved for future use
1	RES	Reserved for future use	R	0	Reserved for future use
0	$\overline{\text{RDY}}/\text{BSY}$	Ready/Busy Status	R	0	Device is ready
				1	Device is busy with an internal operation

1. Only bits 4 and 3 of Status Register Byte 2 are modified when using the Write Status Register Byte 2 command.
2. R/W = Readable and Writeable; R = Readable only.



**Figure 18. Read Status Register**

## 11.2 Write Status Register

The Write Status Register command is used to modify the BPL bit and the BP0 bit of the Status Register. Before the Write Status Register command can be issued, the Write Enable command must have been previously issued to set the WEL bit in the Status Register to a logical 1.

To issue the Write Status Register command, the  $\overline{CS}$  pin must first be asserted and the opcode of 01h must be clocked into the device followed by one byte of data. The one byte of data consists of the BPL bit value, four don't care bits, the BP0 bit value, and two additional don't care bits (see Table 8). Any additional data bytes sent to the device are ignored. When the  $\overline{CS}$  pin is deasserted, the BPL bit and the BP0 bit in the Status Register are modified, and the WEL bit in the Status Register is reset to a logical 0. The value of BP0 and the state of the BPL bit and the  $\overline{WP}$  pin before the Write Status Register command was executed (the prior state of the BPL bit and the state of the  $\overline{WP}$  pin when the  $\overline{CS}$  pin is deasserted) determines whether or not software protection is changed. See Section 9.4, Protected States and the Write Protect Pin, for more details.

The complete one byte of data must be clocked into the device before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on even byte boundaries (multiples of eight bits); otherwise, the device stops the operation, the state of the BPL and BP0 bits does not change, memory protection status does not change, and the WEL bit in the Status Register is reset to the logical 0 state.

If the  $\overline{WP}$  pin is asserted, then the BPL bit can only be set to a logical 1. If an attempt is made to reset the BPL bit to a logical 0 while the  $\overline{WP}$  pin is asserted, then the Write Status Register Byte command is ignored, and the WEL bit in the Status Register is reset to the logical 0 state. In order to reset the BPL bit to a logical 0, the  $\overline{WP}$  pin must be deasserted.

Table 8. Write Status Register Format

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
BPL	X	X	X	X	BP0	X	X

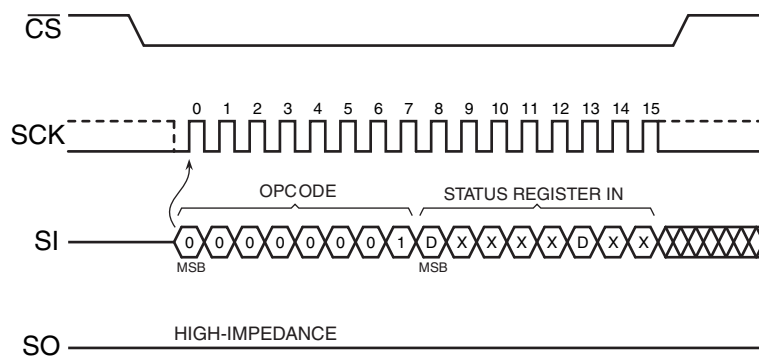


Figure 19. Write Status Register

### 11.3 Write Status Register Byte 2

The Write Status Register Byte 2 command is used to modify the RSTE. Using the Write Status Register Byte 2 command is the only way to modify the RSTE in the Status Register during normal device operation. Before the Write Status Register Byte 2 command can be issued, the Write Enable command must have been previously issued to set the WEL bit in the Status Register to a Logical 1.

To issue the Write Status Register Byte 2 command, the  $\overline{CS}$  pin must first be asserted and then the opcode 31h must be clocked into the device followed by one byte of data. The one byte of data consists of three don't-care bits, the RSTE bit value, and four additional don't-care bits (see Table 9). Any additional data bytes sent to the device are ignored. When the  $\overline{CS}$  pin is deasserted, the RSTE bit in the Status Register is modified, and the WEL bit in the Status Register is reset to a Logical 0.

The complete one byte of data must be clocked into the device before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on even byte boundaries (multiples of eight bits); otherwise, the device stops the operation, the state of the RSTE bit does not change, and the WEL bit in the Status Register is reset to the Logical 0 state.

Table 9. Write Status Register Byte 2 Format

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
X	X	X	RSTE	X	X	X	X

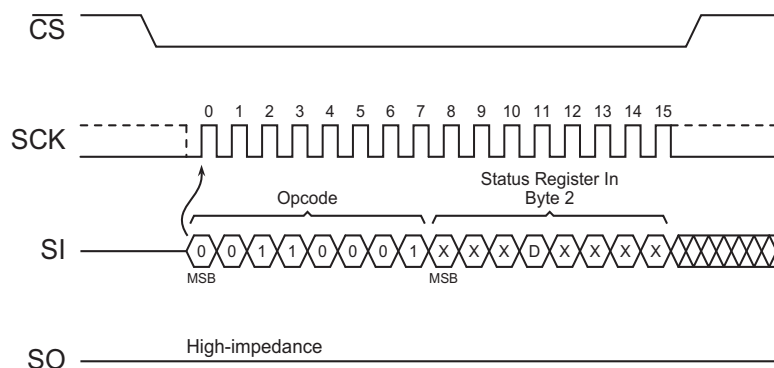


Figure 20. Write Status Register Byte 2

## 12. Other Commands and Functions

### 12.1 Read Manufacturer and Device ID

Identification information can be read from the device to enable systems to electronically query and identify the device while it is in system. The identification method and the command opcode comply with the JEDEC standard for *Manufacturer and Device ID Read Methodology for SPI Compatible Serial Interface Memory Devices*. The type of information that can be read from the device includes the JEDEC defined Manufacturer ID, the vendor specific Device ID, and the vendor specific Extended Device Information.

Since not all Flash devices are capable of operating at very high clock frequencies, applications must be designed to read the identification information from the devices at a reasonably low clock frequency to ensure all devices used in the application can be identified properly. Once the identification process is complete, the application can increase the clock frequency to accommodate specific Flash devices that are capable of operating at the higher clock frequencies.

To read the identification information, the  $\overline{CS}$  pin must first be asserted and the opcode of 9Fh must be clocked into the device. After the opcode has been clocked in, the device begins outputting the identification data on the SO pin during the subsequent clock cycles. The first byte that is output is the Manufacturer ID, followed by two bytes of Device ID information. The fourth byte output is the Extended Device Information String Length, which is 00h, indicating that no Extended Device Information follows. After the Extended Device Information String Length byte is output, the SO pin goes into a high-impedance state; thus, additional clock cycles have no effect on the SO pin and no data are output. As indicated in the JEDEC standard, reading the Extended Device Information String Length and any subsequent data is optional. Deasserting the  $\overline{CS}$  pin stops the Manufacturer and Device ID read operation and puts the SO pin into a high-impedance state. The  $\overline{CS}$  pin can be deasserted at any time and does not require that a full byte of data be read.

**Table 10. Manufacturer and Device ID Information**

Byte No.	Data Type	Value
1	Manufacturer ID	1Fh
2	Device ID (Part 1)	42h
3	Device ID (Part 2)	00h
4	Extended Device Information String Length	00h

**Table 11. Manufacturer and Device ID Details**

Data Type	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Hex Value	Details
Manufacturer ID	JEDEC Assigned Code								1Fh	JEDEC Code: 0001 1111 (1Fh for Renesas Electronics)
	0	0	0	1	1	1	1	1		
Device ID (Part 1)	Family Code			Density Code					42h	Family Code: 010 Density Code: 00010 (1-Mbit)
	0	1	0	0	0	0	1	0		
Device ID (Part 2)	Sub Code			Product Version Code					00h	Sub Code 000 (Standard series) Product Version: 00000
	0	0	0	0	0	0	0	0		

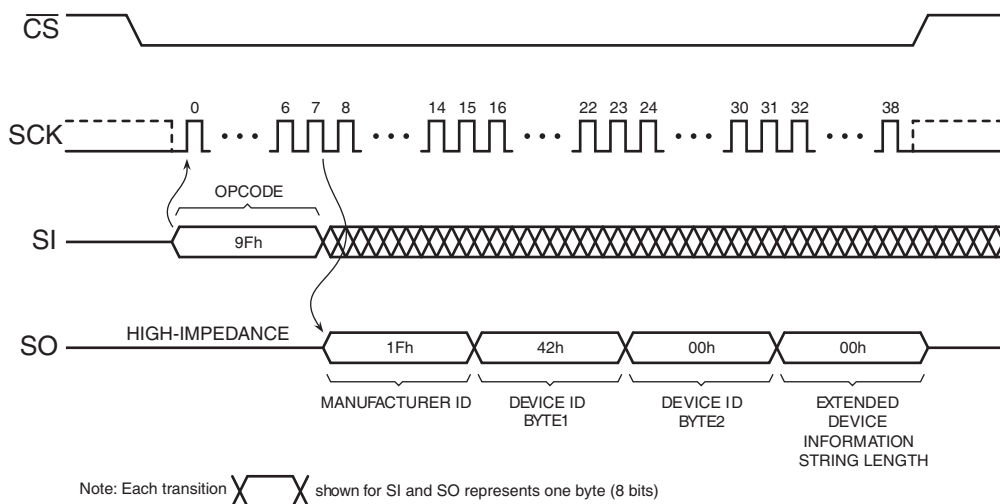


Figure 21. Read Manufacturer and Device ID

## 12.2 Read ID (Legacy Command)

Identification information can be read from the device to enable systems to electronically query and identify the device while it is in system. The preferred method for doing so is the JEDEC standard method described in [Section 12.1](#); however, the legacy Read ID command is supported on the AT25DN011 to enable backwards compatibility to previous generation devices.

To read the identification information, the  $\overline{CS}$  pin must first be asserted and the opcode of 15h must be clocked into the device. After the opcode has been clocked in, the device begins outputting the identification data on the SO pin during the subsequent clock cycles. The first byte output is the Manufacturer ID of 1Fh followed by a single byte of data representing a device code of 65h. After the device code is output, the SO pin goes into a high-impedance state; thus, additional clock cycles have no effect on the SO pin, and no data are output.

Deasserting the  $\overline{CS}$  pin stops the Read ID operation and puts the SO pin into a high-impedance state. The  $\overline{CS}$  pin can be deasserted at any time and does not require that a full byte of data read.

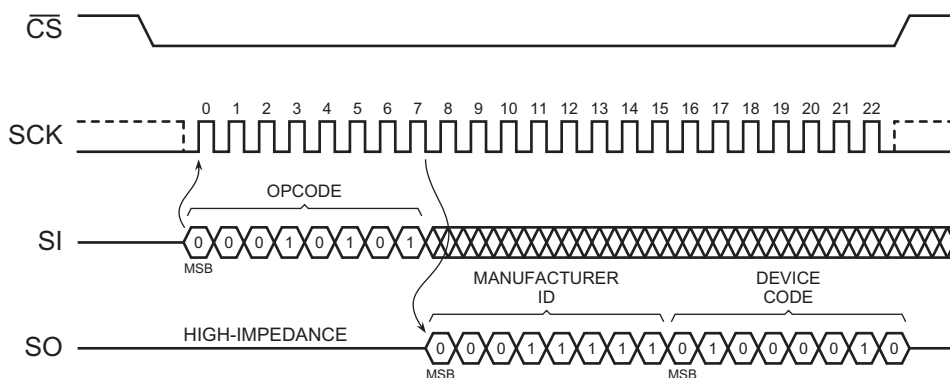


Figure 22. Read ID (Legacy Command)

### 12.3 Deep Power-Down

During normal operation, the device is placed in the standby mode to consume less power as long as the  $\overline{CS}$  pin remains deasserted and no internal operation is in progress. The Deep Power-Down command offers the ability to place the device into an even lower power consumption state called the Deep Power-Down mode.

When the device is in the Deep Power-Down mode, all commands, including the Read Status Register command, are ignored, with the exception of the Resume from Deep Power-Down command. Since all commands are ignored, the mode can be used as an extra protection mechanism against program and erase operations.

Entering the Deep Power-Down mode is done by simply asserting the  $\overline{CS}$  pin, clocking in the opcode of B9h, and then deasserting the  $\overline{CS}$  pin. Any additional data clocked into the device after the opcode are ignored. When the  $\overline{CS}$  pin is deasserted, the device enters the Deep Power-Down mode within the maximum time of  $t_{EDPD}$ .

The complete opcode must be clocked in before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, the device stops the operation and returns to the standby mode once the  $\overline{CS}$  pin is deasserted. Also, the device defaults to the standby mode after a power-cycle.

The Deep Power-Down command is ignored if an internally self-timed operation such as a program or erase cycle is in progress. The Deep Power-Down command must be reissued after the internally self-timed operation has been completed in order for the device to enter the Deep Power-Down mode.

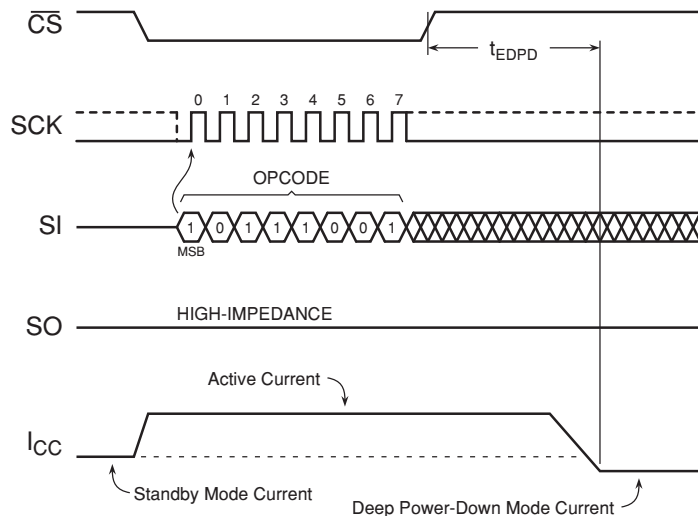


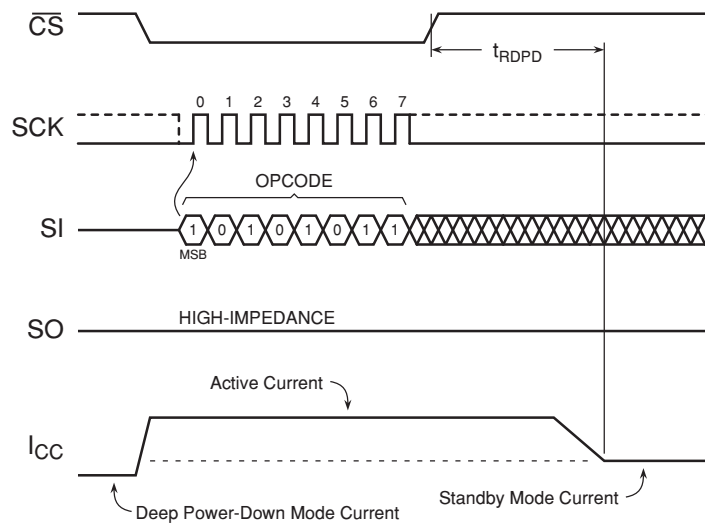
Figure 23. Deep Power-Down

## 12.4 Resume from Deep Power-Down

In order to exit the Deep Power-Down mode and resume normal device operation, the Resume from Deep Power-Down command must be issued. The Resume from Deep Power-Down command is the only command that the device recognizes while in the Deep Power-Down mode.

To resume from the Deep Power-Down mode, the  $\overline{CS}$  pin must first be asserted and opcode of ABh must be clocked into the device. Any additional data clocked into the device after the opcode are ignored. When the  $\overline{CS}$  pin is deasserted, the device exits the Deep Power-Down mode within the maximum time of  $t_{RDPD}$  and returns to the standby mode. After the device has returned to the standby mode, normal command operations such as Read Array can be resumed.

If the complete opcode is not clocked in before the  $\overline{CS}$  pin is deasserted, or if the  $\overline{CS}$  pin is not deasserted on an even byte boundary (multiples of eight bits), then the device stops the operation and returns to the Deep Power-Down mode.



**Figure 24. Resume from Deep Power-Down**

## 12.5 Ultra-Deep Power-Down

The Ultra-Deep Power-Down mode allows the device to further reduce its energy consumption compared to the existing standby and Deep Power-Down modes by shutting down additional internal circuitry. When the device is in the Ultra-Deep Power-Down mode, all commands including the Status Register Read and Resume from Deep Power-Down commands are ignored. Since all commands are ignored, the mode can be used as an extra protection mechanism against inadvertent or unintentional program and erase operations. Entering the Ultra-Deep Power-Down mode is done by simply asserting the  $\overline{CS}$  pin, clocking in the opcode 79h, and then deasserting the  $\overline{CS}$  pin. Any additional data clocked into the device after the opcode are ignored. When the  $\overline{CS}$  pin is deasserted, the device enters the Ultra-Deep Power-Down mode within the maximum time of  $t_{EUDPD}$ .

The complete opcode must be clocked in before the  $\overline{CS}$  pin is deasserted; otherwise, the device stops the operation and returns to the standby mode once the  $\overline{CS}$  pin is deasserted. Also, the device defaults to the standby mode after a power cycle. The Ultra-Deep Power-Down command is ignored if an internally self-timed operation such as a program or erase cycle is in progress.

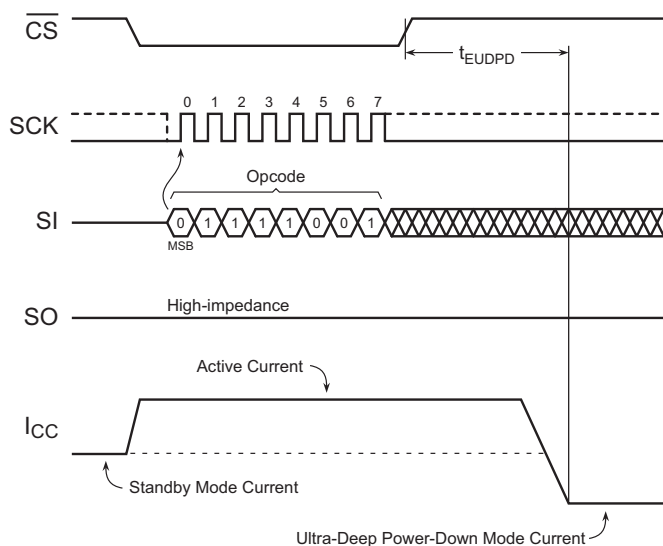


Figure 25. Ultra-Deep Power-Down



## 12.6 Exit Ultra-Deep Power-Down

To exit from the Ultra-Deep Power-Down mode, any one of three operations can be performed, which are described in the following subsections.

### 12.6.1 Chip Select Toggle

The  $\overline{CS}$  pin must simply be pulsed by asserting the  $\overline{CS}$  pin, waiting the minimum necessary  $t_{CSLU}$  time, and then deasserting the  $\overline{CS}$  pin again. To facilitate simple software development, a dummy byte opcode can also be entered while the  $\overline{CS}$  pin is being pulsed; the dummy byte opcode is simply ignored by the device in this case. After the  $\overline{CS}$  pin has been deasserted, the device exits from the Ultra-Deep Power-Down mode and returns to the standby mode within a maximum time of  $t_{XUDPD}$ . If the  $\overline{CS}$  pin is reasserted before the  $t_{XUDPD}$  time has elapsed in an attempt to start a new operation, then that operation is ignored, and nothing is done.

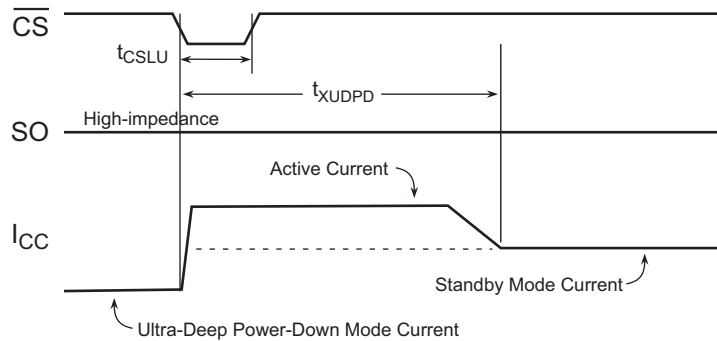


Figure 26. Exit Ultra-Deep Power-Down (Chip Select Toggle)

### 12.6.2 Chip Select Low

By asserting the  $\overline{CS}$  pin, waiting the minimum necessary  $t_{XUDPD}$  time, and then clocking in the first bit of the next Opcode command cycle. If the first bit of the next command is clocked in before the  $t_{XUDPD}$  time has elapsed, the device exits Ultra-Deep Power-Down; however, the intended operation is ignored.

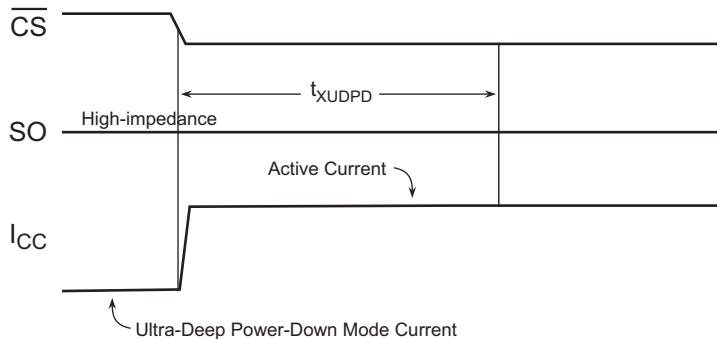


Figure 27. Exit Ultra-Deep Power-Down (Chip Select Low)

### 12.6.3 Power Cycling

The device can also exit the Ultra-Deep Power Mode by power cycling the device. The system must wait for the device to return to the standby mode before normal command operations can be resumed. Upon recovery from Ultra-Deep Power-Down, all internal registers are at their Power-On default state.

## 12.7 Hold

The  $\overline{\text{HOLD}}$  pin is used to pause the serial communication with the device without having to stop or reset the clock sequence. The Hold mode, however, does not have an effect on any internally self-timed operations such as a program or erase cycle. Thus, if an erase cycle is in progress, asserting the  $\overline{\text{HOLD}}$  pin does not pause the operation, and the erase cycle continues until it is finished.

The Hold mode can only be entered while the  $\overline{\text{CS}}$  pin is asserted. The Hold mode is activated simply by asserting the  $\overline{\text{HOLD}}$  pin during the SCK low pulse. If the  $\overline{\text{HOLD}}$  pin is asserted during the SCK high pulse, then the Hold mode won't be started until the beginning of the next SCK low pulse. The device remains in the Hold mode as long as the  $\overline{\text{HOLD}}$  pin and  $\overline{\text{CS}}$  pin are asserted.

While in the Hold mode, the SO pin is in a high-impedance state. Also, both the SI pin and the SCK pin are ignored. The  $\overline{\text{WP}}$  pin, however, can still be asserted or deasserted while in the Hold mode.

To end the Hold mode and resume serial communication, the  $\overline{\text{HOLD}}$  pin must be deasserted during the SCK low pulse. If the  $\overline{\text{HOLD}}$  pin is deasserted during the SCK high pulse, then the Hold mode won't end until the beginning of the next SCK low pulse.

If the  $\overline{\text{CS}}$  pin is deasserted while the  $\overline{\text{HOLD}}$  pin is still asserted, then any operation that was started is stopped, and the device resets the WEL bit in the Status Register to the logical 0 state.

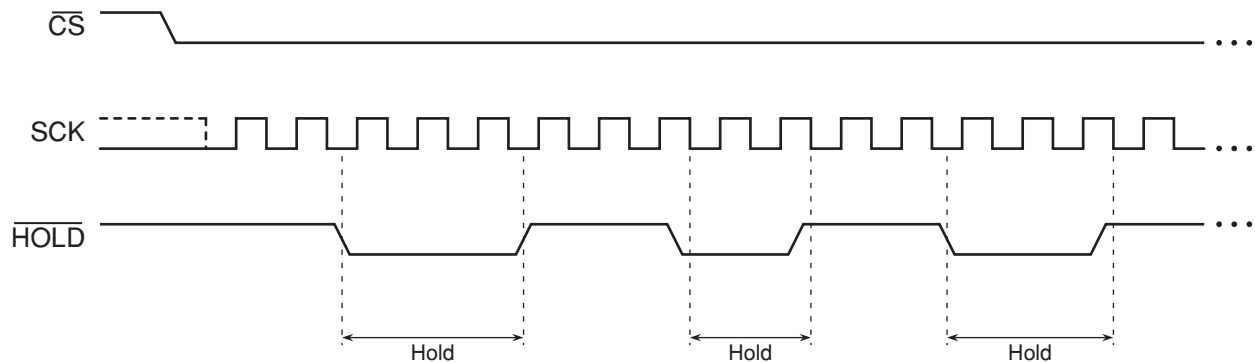


Figure 28. Hold Mode

## 12.8 Reset

In some applications, it might be necessary to prematurely stop a program or erase operation rather than wait the hundreds of microseconds or milliseconds necessary for the program or erase operation to complete normally. The Reset command allows a program or erase operation in progress to be ended abruptly and returns the device to an idle state. Since the need to reset the device is immediate, the Write Enable command does not need to be issued prior to the Reset command. Thus, the Reset command operates independently of the state of the WEL bit in the Status Register.

The Reset command can be executed only if the command has been enabled by setting the Reset Enabled (RSTE) bit in the Status Register to a Logical 1 using write status register byte 2 command 31h. Enter this command before a program command is entered. If the Reset command has not been enabled (the RSTE bit is in the Logical 0 state), then any attempts at executing the Reset command are ignored.

To perform a Reset, the  $\overline{CS}$  pin must first be asserted, and then the opcode F0h must be clocked into the device. No address bytes need to be clocked in, but a confirmation byte of D0h must be clocked into the device immediately after the opcode. Any additional data clocked into the device after the confirmation byte are ignored. When the  $\overline{CS}$  pin is deasserted, the program operation currently in progress is stopped within a time of  $t_{SWRST}$ . Since the program or erase operation might not complete before the device is reset, the contents of the page being programmed or erased cannot be guaranteed to be valid.

The Reset command has no effect on the states of the Configuration Register or RSTE bit in the Status Register. The WEL bit, however, is reset to its default state.

The complete opcode and confirmation byte must be clocked into the device before the  $\overline{CS}$  pin is deasserted, and the  $\overline{CS}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, no Reset operation is done.

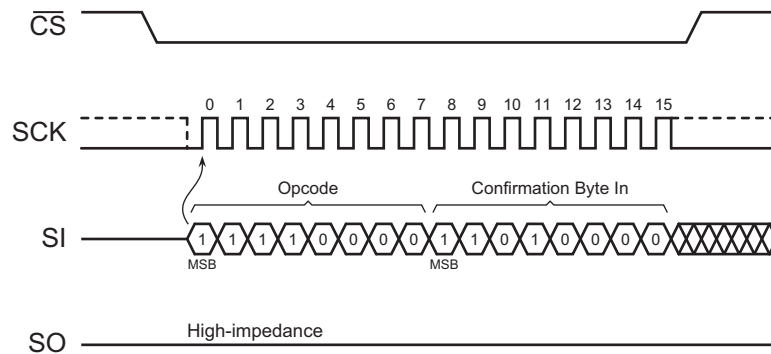


Figure 29. Reset

## 13. Electrical Specifications

### 13.1 Absolute Maximum Ratings

Temperature under Bias	-55 °C to +125 °C
Storage Temperature	-65 °C to +150 °C
All Input Voltages (including NC Pins) with Respect to Ground	-0.6 V to 4.1 V
All Output Voltages with Respect to Ground	-0.6 V to $V_{CC} + 0.5$ V

\*Notice: Stresses beyond those listed for Absolute Maximum can cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions beyond those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods can affect device reliability.

### 13.2 DC and AC Operating Range

Parameter	Comment	Range
Operating Temperature (Case)	Industrial	-40 °C to +85 °C
$V_{CC}$ Power Supply		2.3 V to 3.6 V

### 13.3 DC Characteristics

Symbol	Parameter	Condition	2.3 V to 3.6 V			Units
			Min	Typ	Max	
$I_{UDPD}$	Ultra-Deep Power-Down Current	$\overline{CS} = V_{CC}$ . All other inputs at 0 V or $V_{CC}$		0.35	0.6	$\mu$ A
$I_{DPD}$	Deep Power-Down Current	$\overline{CS} = V_{CC}$ . All other inputs at 0 V or $V_{CC}$		7.5	15	$\mu$ A
$I_{SB}$	Standby Current	$\overline{CS} = V_{CC}$ . All other inputs at 0 V or $V_{CC}$		25	40	$\mu$ A
$I_{CC1}^1$	Active Current, Low Power Read (03h, 0Bh) Operation	$f = 1$ MHz; $I_{OUT} = 0$ mA		6	11	mA
		$f = 20$ MHz; $I_{OUT} = 0$ mA		7	11	mA
$I_{CC2}^1$	Active Current, Read Operation	$f = 50$ MHz; $I_{OUT} = 0$ mA		7	11	mA
		$f = 85$ MHz; $I_{OUT} = 0$ mA		7	12	mA
$I_{CC3}^1$	Active Current, Program Operation	$\overline{CS} = V_{CC}$		12	16	mA
$I_{CC4}^1$	Active Current, Erase Operation	$\overline{CS} = V_{CC}$		12	18	mA
$I_{LI}$	Input Load Current	All inputs at CMOS levels			1	$\mu$ A
$I_{LO}$	Output Leakage Current	All inputs at CMOS levels			1	$\mu$ A
$V_{IL}$	Input Low Voltage				$V_{CC} \times 0.3$	V
$V_{IH}$	Input High Voltage		$V_{CC} \times 0.7$			V
$V_{OL}$	Output Low Voltage	$I_{OL} = 100$ $\mu$ A			0.4	V
$V_{OH}$	Output High Voltage	$I_{OH} = -100$ $\mu$ A	$V_{CC} - 0.2$ V			V

1. Typical values measured at 3.0 V @ 25 °C for the 2.3 V to 3.6 V range.

## 13.4 AC Characteristics

Symbol	Parameter	2.3 V to 3.6 V			Units
		Min	Typ	Max	
$f_{CLK}$	Maximum Clock Frequency for All Operations except 03h and 3Bh Opcodes			104	MHz
$f_{RDLF}$	Maximum Clock Frequency for 03h Opcode (Read Array – Low Frequency)			33	MHz
$f_{RDDO}$	Maximum Clock Frequency for 3B Opcode			50	MHz
$t_{CLKH}$	Clock High Time	4			ns
$t_{CLKL}$	Clock Low Time	4			ns
$t_{CLKR}^1$	Clock Rise Time, Peak-to-Peak (Slew Rate)	0.1			V/ns
$t_{CLKF}^1$	Clock Fall Time, Peak-to-Peak (Slew Rate)	0.1			V/ns
$t_{CSH}$	Chip Select High Time	25			ns
$t_{CSLS}$	Chip Select Low Setup Time (relative to Clock)	5			ns
$t_{CSLH}$	Chip Select Low Hold Time (relative to Clock)	5			ns
$t_{CSHS}$	Chip Select High Setup Time (relative to Clock)	5			ns
$t_{CSHH}$	Chip Select High Hold Time (relative to Clock)	5			ns
$t_{DS}$	Data In Setup Time	2			ns
$t_{DH}$	Data In Hold Time	1			ns
$t_{DIS}^1$	Output Disable Time			6	ns
$t_V$	Output Valid Time			6	ns
$t_{OH}$	Output Hold Time	0			ns
$t_{HLS}$	$\overline{HOLD}$ Low Setup Time (relative to Clock)	5			ns
$t_{HLH}$	$\overline{HOLD}$ Low Hold Time (relative to Clock)	5			ns
$t_{HHS}$	$\overline{HOLD}$ High Setup Time (relative to Clock)	5			ns
$t_{HHH}$	$\overline{HOLD}$ High Hold Time (relative to Clock)	5			ns
$t_{HLQZ}^1$	$\overline{HOLD}$ Low to Output High-Z			6	ns
$t_{HHQX}^1$	$\overline{HOLD}$ High to Output Low-Z			6	ns
$t_{WPS}^{1,2}$	Write Protect Setup Time	20			ns
$t_{WPH}^{1,2}$	Write Protect Hold Time	100			ns
$t_{EDPD}^1$	Chip Select High to Deep Power-Down			2	$\mu$ s
$t_{EUDPD}$	Chip Select High to Ultra-Deep Power-Down			3	$\mu$ s
$t_{SWRST}$	Software Reset Time			50	$\mu$ s
$t_{CSLU}$	Minimum Chip Select Low to Exit Ultra-Deep Power-Down	20			ns
$t_{XUDPD}$	Exit Ultra-Deep Power-Down Time	70			$\mu$ s
$t_{RDPD}^1$	Chip Select High to Standby Mode			8	$\mu$ s

1. Not 100% tested (value guaranteed by design and characterization).

2. Only applicable as a constraint for the Write Status Register command when BPL = 1.

### 13.5 Program and Erase Characteristics

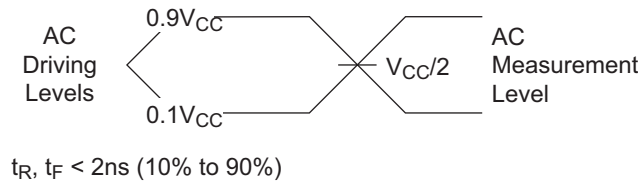
Symbol	Parameter		2.3 V - 3.6 V			
			Min	Typ	Max	Units
$t_{PP}^1$	Page Program Time	256 Bytes		1.25	1.75	ms
$t_{BP}$	Byte Program Time			8		$\mu$ s
$t_{PE}$	Page Erase Time	256 Bytes		6	20	ms
$t_{BLKE}^1$	Block Erase Time	4 kBytes		35	50	ms
		32 kBytes		250	350	
$t_{CHPE}^{1,2}$	Chip Erase Time			1000	1400	ms
$t_{OTPP}^1$	OTP Security Register Program Time			400	950	$\mu$ s
$t_{WRSR}^2$	Write Status Register Time			20	40	ms

1. Maximum values indicate worst-case performance after 100,000 erase/program cycles.
2. Not 100% tested (value guaranteed by design and characterization).

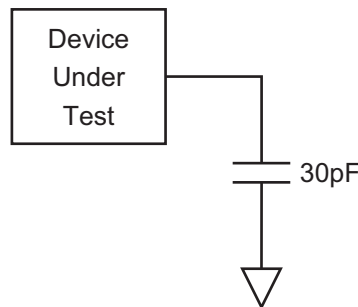
### 13.6 Power-Up Conditions

Symbol	Parameter	Min	Max	Units
$f_{VCSL}$	Minimum $V_{CC}$ to Chip Select Low Time	70		$\mu$ s
$f_{PUW}$	Power-up Device Delay Before Program or Erase Allowed		5	ms
$f_{POR}$	Power-on Reset Voltage	1.6	22	V

### 13.7 Input Test Waveforms and Measurement Levels



### 13.8 Output Test Load



## 14. AC Waveforms

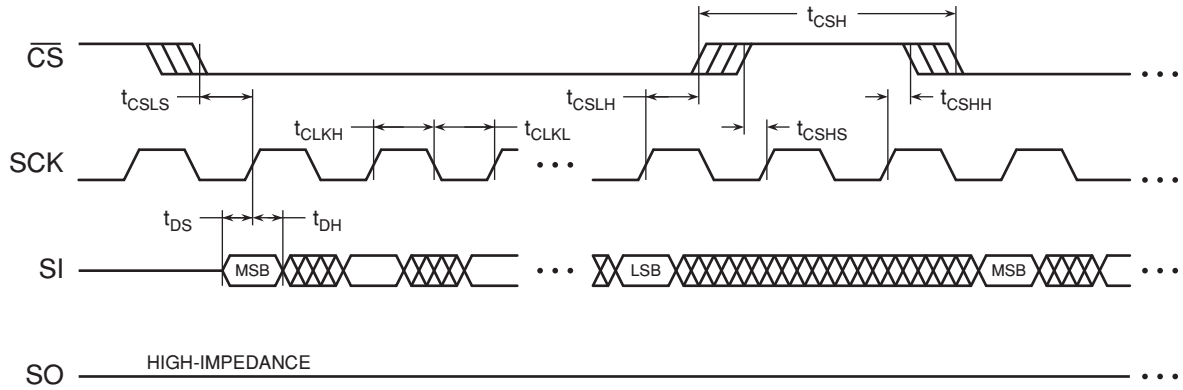


Figure 30. Serial Input Timing

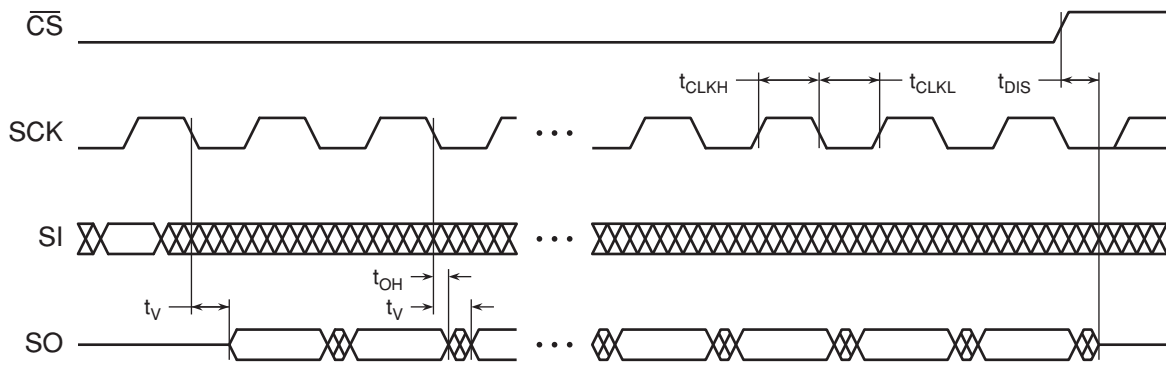


Figure 31. Serial Output Timing

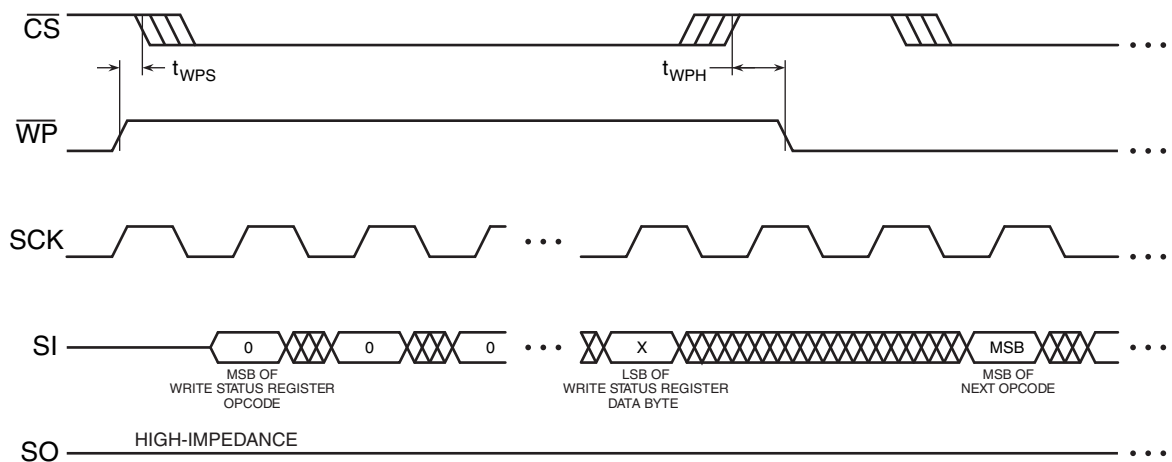


Figure 32.  $\overline{WP}$  Timing for Write Status Register Command When BPL = 1

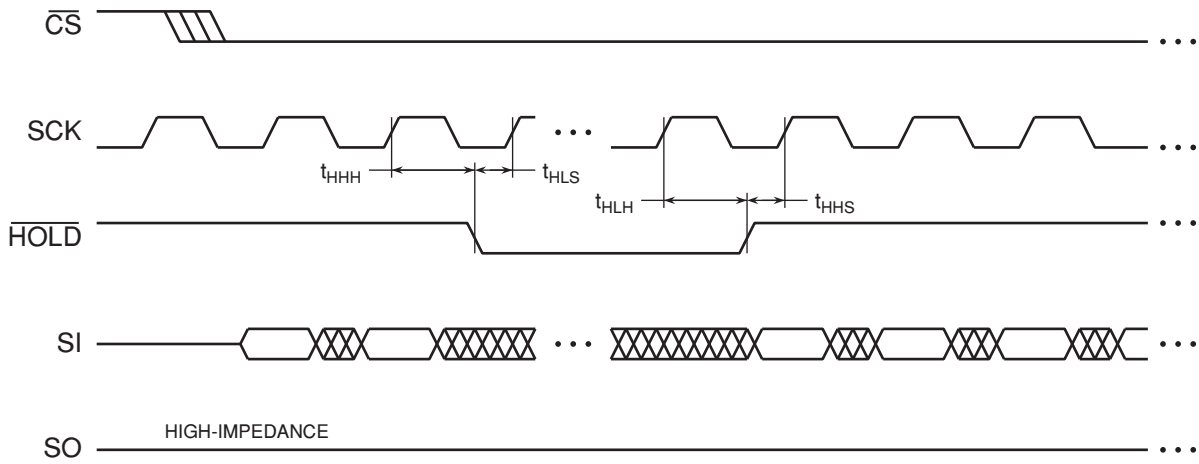


Figure 33.  $\overline{HOLD}$  Timing - Serial Input

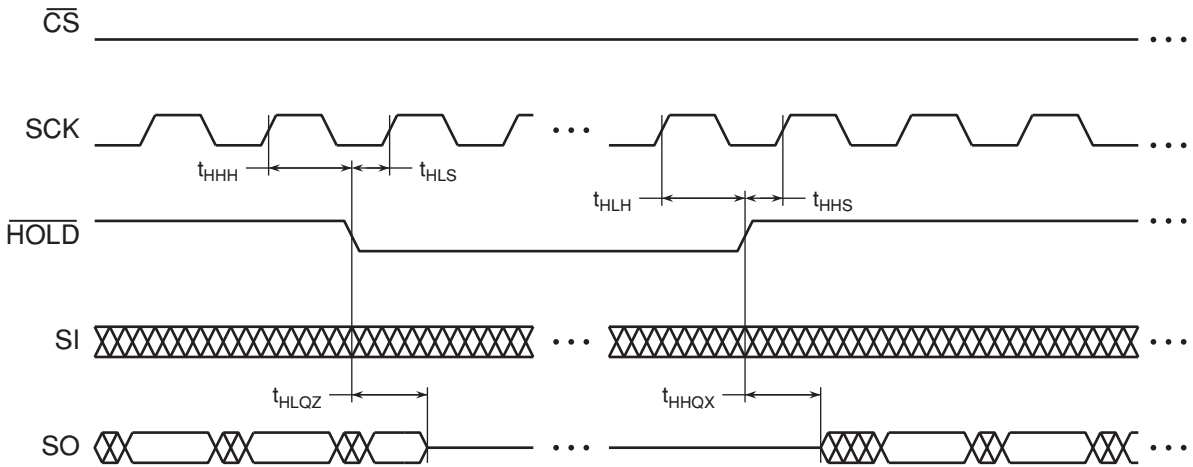
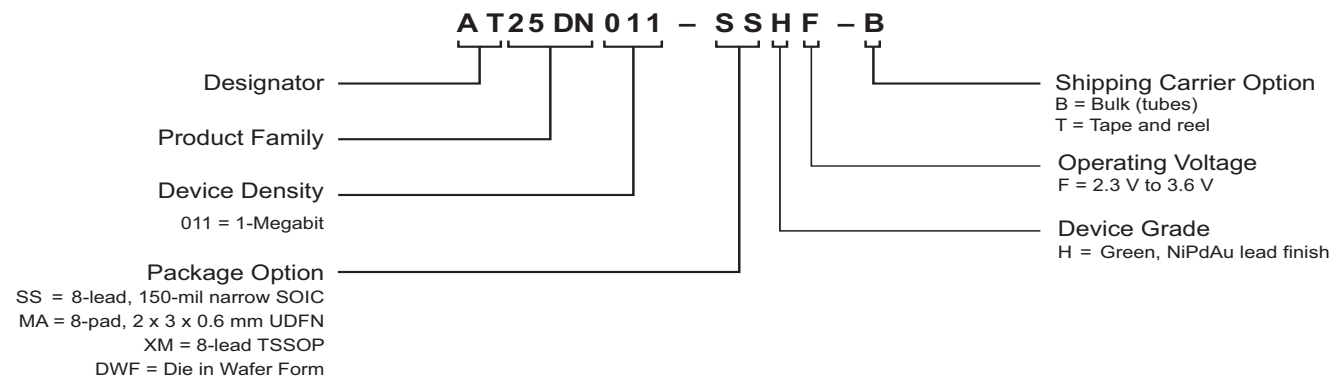


Figure 34.  $\overline{HOLD}$  Timing - Serial Output



## 15. Ordering Information

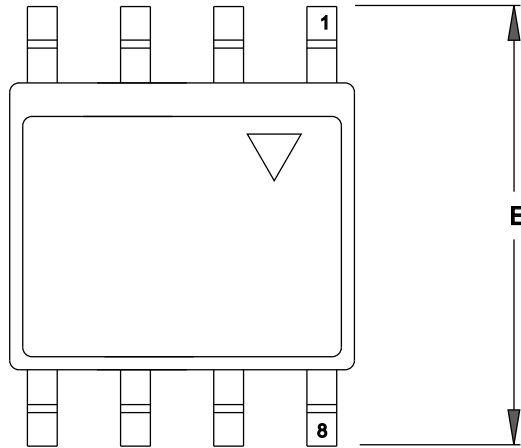


Ordering Code <sup>1</sup>	Package	Lead Finish	Operating Voltage	Max. Freq. (MHz)	Temperature Range
AT25DN011-SSHFB AT25DN011-SSHFT	8-lead, 150-mil narrow, Plastic Gull Wing Small Outline Package (JEDEC SOIC)	NiPdAu	2.3 V to 3.6 V	104	Industrial (-40 °C to +85 °C)
AT25DN011-MAHFT	8-pad, 2 x 3 x 0.6 mm, Thermally Enhanced Plastic Ultra Thin Dual Flat No Lead Package (UDFN)				
AT25DN011-XMHFB AT25DN011-XMHFT	8-lead, Thin Small Outline Package				
AT25DN011-DWF	Die in Wafer Form	N/A			

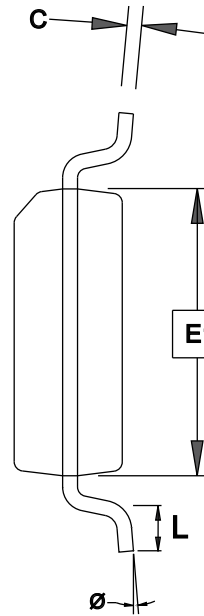
1. The shipping carrier option code is not marked on the devices.

## 16. Packaging Information

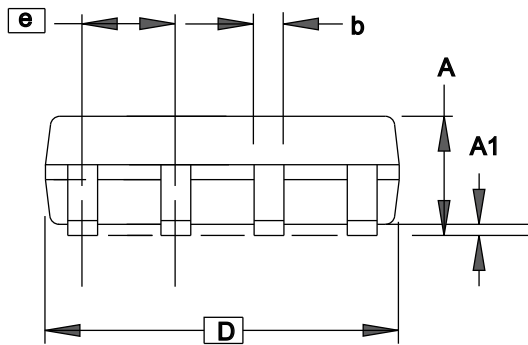
### 16.1 8-Lead, 150-mil Narrow JEDEC SOIC



TOP VIEW



END VIEW



SIDE VIEW

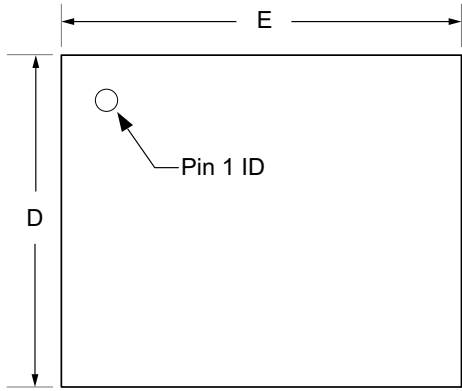
Notes: This drawing is for general information only.  
See JEDEC Drawing MS-012, Variation AA,  
for proper dimensions, tolerances, datums, etc.

COMMON DIMENSIONS  
(Unit of Measure = mm)

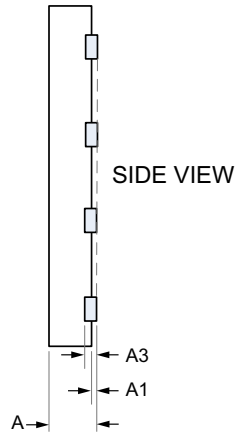
SYMBOL	MIN	NOM	MAX	NOTE
A	1.35	1.60	1.75	
A1	0.10	0.18	0.25	
b	0.31	0.43	0.51	
C	0.17	0.22	0.25	
D	4.80	4.83	5.05	
E1	3.81	3.90	3.99	
E	5.79	6.00	6.20	
e	1.27 BSC			
L	0.40	0.60	1.27	
Ø	0°	3.75	8°	



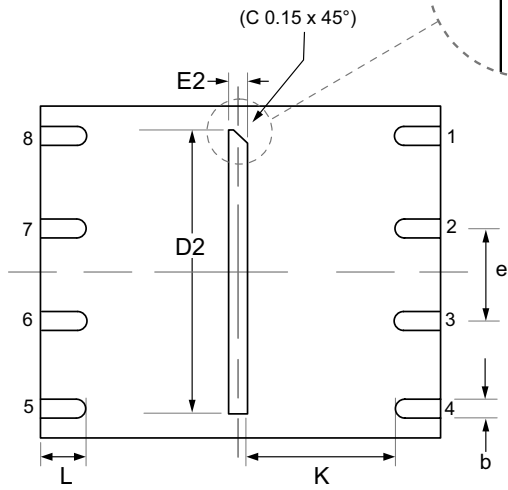
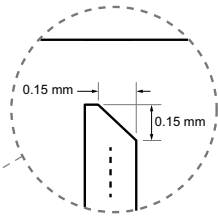
### 16.2 2 x 3 x 0.6 mm UDFN



TOP VIEW



SIDE VIEW



BOTTOM VIEW

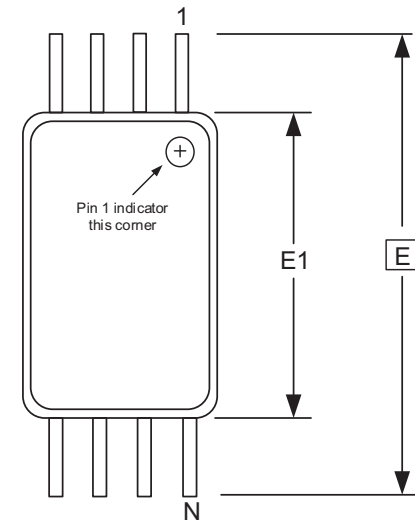
COMMON DIMENSIONS  
(Unit of Measure = mm)

SYMBOL	MIN	NOM	MAX
A	0.50	0.55	0.60
A1	0.00	0.02	0.05
A3	0.150 REF		
b	0.20	0.25	0.30
D	2.00 BSC		
D2	1.55	1.60	1.65
E	3.00 BSC		
e	0.50 BSC		
E2	0.15	0.20	0.25
K	0.20		
L	0.40	0.45	0.50

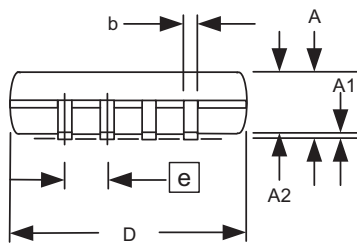
- Notes: 1. This package conforms to JEDEC reference MO-229, Saw Singulation.  
 2. The terminal #1 ID is a Laser-marked feature.



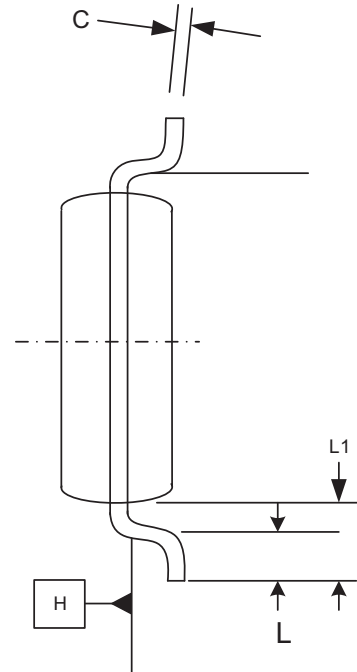
### 16.3 TSSOP



TOP VIEW



SIDE VIEW



END VIEW

COMMON DIMENSIONS  
(Unit of Measure = mm)

SYMBOL	MIN	NOM	MAX	NOTE
A	—	—	1.20	
A1	0.05	—	0.15	
A2	0.80	1.00	1.05	
D	2.90	3.00	3.10	2, 5
E	6.40 BSC			
E1	4.30	4.40	4.50	3, 5
b	0.19	—	0.30	4
e	0.65 BSC			
L	0.45	0.60	0.75	
L1	1.00 REF			
C	0.09	—	0.20	

Notes:

1. This drawing is for general information only. See JEDEC Drawing MO-153, Variation AA, for proper dimensions, tolerances, datums, etc.
2. Dimension D does not include mold Flash, protrusions, or gate burrs. Mold Flash, protrusions, and gate burrs shall not exceed 0.15 mm (0.006 in) per side.
3. Dimension E1 does not include inter-lead Flash or protrusions, Inter-lead Flash and protrusions shall not exceed 0.25 mm (0.01 in) per side.
4. Dimension b does not include Dambar protrusion. Allowable Dambar protrusion shall be 0.08 mm total in excess of the b dimension at maximum material condition. Dambar cannot be located on the lower radius of the foot. Minimum space between protrusion and adjacent lead is 0.07 mm.
5. Dimensions D and E1 to be determined at Datum Plane H.



## 17. Revision History

Revision Level – Release Date	History
A – January 2014	Initial release.
B – May 2014	Removed “All inputs (SCK, CS, WP and HOLD) are guaranteed by design to be 5V tolerant from Table 13-3. Updated AC and DC characteristics. Removed tray carrier option from DFN parts. Changed datasheet status to Preliminary.
C – October 2014	Corrected last byte address in sections 6, 7.1, and 7.2. Corrected 8S1 and 8MA3 package outline drawings. Updated $I_{UDPD}$ description in Table 13.3. Corrected $I_{OL}$ . Corrected product density error on Ordering Detail diagram. Updated Table 9-1 (Protected Memory Address). Updated AC and DC specifications.
D – November 2015	Updated $I_{DPD}$ and $I_{SB}$ specification conditions. Updated document status from Preliminary to Complete.
E – February 2017	Added patent information. Updated description in Section 8.2 (Page Erase).
F – August 2017	Added Die in Wafer Form (DWF) package option.
G – March 2021	Added NOM numbers in the Common Dimensions table in 16.1.
H – August 2021	Updated UDFN package drawing (Section 16.2).
I – December 2021	Updated UDFN package drawing (Section 16.2).
J – February 2022	Changed company logo to Renesas.
K – November 2023	Applied new corporate template to document.

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